

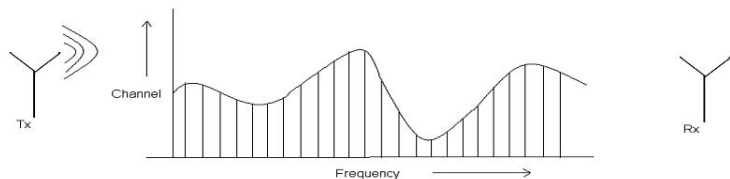
Optimal Resource Allocation for OFDM Uplink Communication: A Primal-Dual Approach

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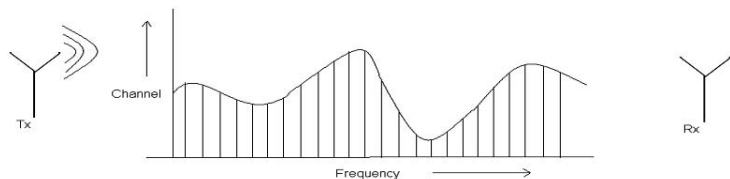
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OFDM Systems



- Frequency band divided into several parallel **orthogonal** carriers/tones.
- High spectrum efficiency.
- Eliminate inter-symbol-interference (ICI) due to multi-path fading.
- Applications: WiMAX (802.16), Wi-Fi (802.11a/g), DSL, etc.

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- Applications: WiMAX (802.16), Wi-Fi (802.11a/g), DSL, etc.
- How to perform **distributed** and **efficient** resource allocation?

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- ▶ Motivated by **channel-aware** and **gradient-based** scheduling
- ▶ **Myopic** policy, requires no knowledge of channel or arrival statistics.

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- Resource allocation = **weighted rate maximization**
 - ▶ “Network”: assigning carriers to users
 - ▶ “User”: assign power over carriers

Achievable Rate Per Carrier

- If a carrier j is allocated to a **single** user i :
 - ▶ $r_{ij} = \log(1 + SNR) = \log(1 + p_{ij}e_{ij})$
 - ▶ p_{ij} = power user i allocates to carrier j .
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- Why “time-sharing”?
 - ▶ **Convexify** the problem.
 - ▶ Can be achieved in practice.
 - ▶ **Not affecting** the optimal objective value with large number of carriers.

OFDMA rate region

- Rate region:

$$\mathcal{R}(\mathbf{e}) = \left\{ \mathbf{r} : r_i = \sum_j x_{ij} \log \left(1 + \frac{p_{ij} e_{ij}}{x_{ij}} \right), \forall i \right\}$$

where $(\mathbf{x}, \mathbf{p}) \in \mathcal{X}$ such that

- ▶ $(\mathbf{x}, \mathbf{p}) \geq \mathbf{0}$
- ▶ Each user i : $\sum_j p_{ij} \leq P_i$ (uplink transmission power constraint)
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- Variation: **sub-channelization** (bundle carriers to reduce overhead)
 - ▶ Interleaved (802.16 standard mode)
 - ▶ Adjacent (Band AMC mode)
 - ▶ Random (e.g. frequency hopped)

Weighted Rate Maximization Problem

$$\max_{\mathbf{r} \in \mathcal{R}(\mathbf{e})} \mathbf{w} \cdot \mathbf{r} = \max_{(\mathbf{x}, \mathbf{p}) \in \mathcal{X}} \sum_i w_i \sum_j x_{ij} \log \left(1 + \frac{p_{ij} e_{ij}}{x_{ij}} \right) \quad (\text{OPT})$$

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- Concave maximization problem
- Non-strictly concave: typically many local/global optimal allocations
- What is known: **centralized** optimal and (very good) heuristic algorithms [HSBA'07]
- What we will show: **distributed** optimal algorithm [This talk]

Centralized Optimal Algorithm [HSBA'07]

- Problem (OPT) is convex and satisfies Slater's condition
⇒ No duality gap.
- Consider **Lagrangian**:

$$L(\mathbf{x}, \mathbf{p}, \boldsymbol{\lambda}, \boldsymbol{\mu}) := \sum_i w_i \sum_j x_{ij} \log \left(1 + \frac{p_{ij} e_{ij}}{x_{ij}} \right) + \sum_i \lambda_i \left(P_i - \sum_j p_{ij} \right) + \sum_j \mu_j \left(1 - \sum_i x_{ij} \right).$$

- Dual function:

$$L(\boldsymbol{\lambda}, \boldsymbol{\mu}) = \max_{(\mathbf{x}, \mathbf{p}) \in \mathcal{X}} L(\mathbf{x}, \mathbf{p}, \boldsymbol{\lambda}, \boldsymbol{\mu})$$

- Optimal objective value to Problem (OPT):

$$V^* = \min_{(\boldsymbol{\lambda}, \boldsymbol{\mu}) \geq \mathbf{0}} L(\boldsymbol{\lambda}, \boldsymbol{\mu})$$

A Dual-Based Centralized Algorithm [HSBA'07]

- 1 For fixed λ , analytically solve for $\mu(\lambda)$, $\mathbf{x}(\lambda)$ and $\mathbf{p}(\lambda)$.
- 2 Multi-dimensional subgradient search of optimal λ^* .
- 3 Find optimal \mathbf{x}^* and \mathbf{p}^* by solving a system of linear equations (multiple solutions).

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Properties:

- **Centralized** computation (step 3)
- **Slow** convergence (step 2)
- **High** signaling overhead (steps 1, 2 and 3): all channel information and power constraints

Distributed Primal-Dual Algorithm

- Try to reach a **saddle point** of Lagrangian $L(\mathbf{x}, \mathbf{p}, \boldsymbol{\lambda}, \boldsymbol{\mu})$:

$$\dot{x}_{ij} = k_{ij}^x \cdot \partial L / \partial x_{ij} = k_{ij}^x (f_{ij}(x_{ij}, p_{ij}) - \mu_j)_{x_{ij}}^+, \text{ (mobile)}$$

$$\dot{p}_{ij} = k_{ij}^p \cdot \partial L / \partial p_{ij} = k_{ij}^p (g_{ij}(x_{ij}, p_{ij}) - \lambda_i)_{p_{ij}}^+, \text{ (mobile)}$$

$$\dot{\lambda}_i = k_i^\lambda \cdot \partial L / \partial \lambda_i = k_i^\lambda \left(\sum_j p_{ij} - P_i \right)_{\lambda_i}^+, \text{ (mobile)}$$

$$\dot{\mu}_j = k_j^\mu \cdot \partial L / \partial \mu_j = k_j^\mu \left(\sum_i x_{ij} - 1 \right)_{\mu_j}^+, \text{ (base station)}$$

Distributed Primal-Dual Algorithm

Advantages:

- **Distributed** and **simple** updates by mobiles and base station
- **Low** communication overhead: user feedback x_{ij} 's and base station announces μ_j

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Challenge:

- How to achieve **global convergence** with **non-strictly concave** objective functions
 - ▶ Difficult in general
 - ▶ Example: multi-path routing [Voice'06]

Convergence Result: Part I

Theorem (Convergence to an Invariant Set)

- All trajectories of the primal-dual system converge to an invariant set V_0 *globally* and *asymptotically*.
- All *optimal* solutions of Problem (OPT) are contained in set V_0 .

Convergence Result: Part I

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- **Proof:** constructing a proper Lyapunov function.
 - **Question:** Will set V_0 contain non-optimal solutions?

Convergence Result: Part II

Theorem (Convergence to a Global Optimal Solution)

All trajectories of the primal-dual system globally converge to optimal solutions of Problem (OPT) under properly chosen stepsizes.

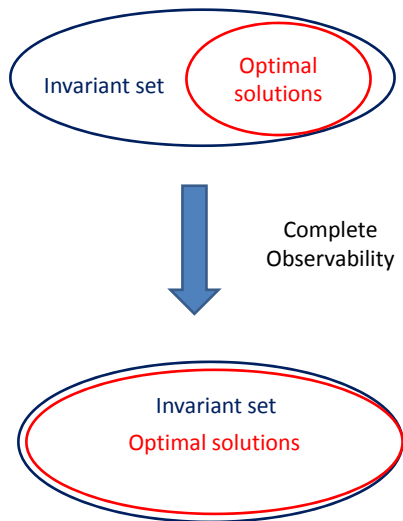
Convergence Result: Part II

Theorem (Convergence to a Global Optimal Solution)

All trajectories of the primal-dual system globally converge to optimal solutions of Problem (OPT) under properly chosen stepsizes.

- Set V_0 only contains optimal solutions.
- Stepsize choice is easy: e.g., not all k_j^μ 's are the same.
- **Proof:**
 - ▶ Over set V_0 , the nonlinear system reduces to a **linear** one.
 - ▶ The linear system is **marginal stable**.
 - ▶ Set V_0 contains only the optimal solution if (λ, μ) is **completely observable** from $B[x^T, p^T]^T$.

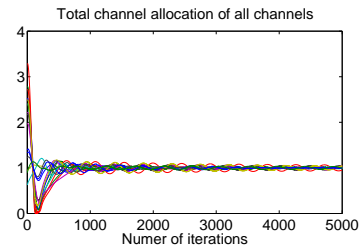
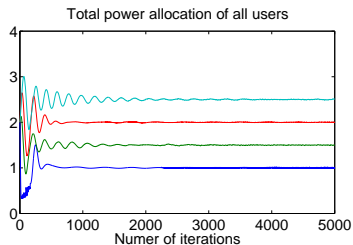
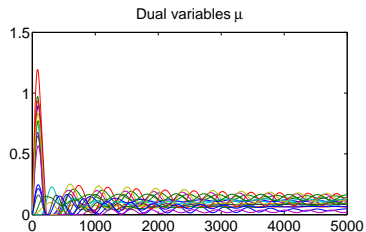
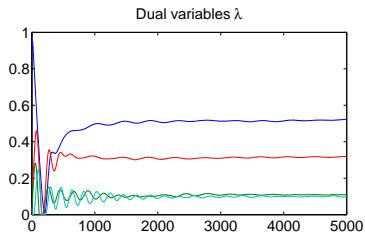
Relationship



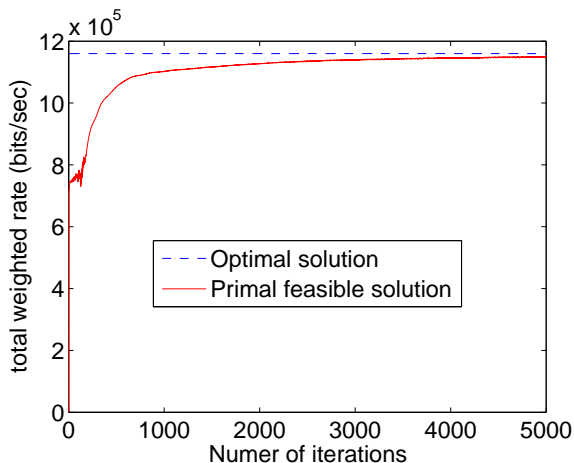
Simulation Set-up

- Single cell, $M = 4$ users.
- $e_{ij} = (\text{fixed location-based term}) \times (\text{frequency selective fast fading})$
 - ▶ Fixed term = empirical distribution.
 - ▶ frequency selective term = block fading in time (2msec coh. time $\approx 250\text{Hz}$ Doppler); standard ref. mobile delay spread (6 taps, $1 \mu\text{sec}$).
- 5 MHz BW, 512 carriers.
- Adjacent channelization, 8 carriers/subchannel.
- Resource allocation over 20 OFDM symbols.
- Randomly generated weights in $[0, 1]$.

Convergence of Primal and Dual Variables



Convergence of Total Weighted Rate



Achieve 90% within 500 iterations

Conclusions

- **Topic:** optimal resource allocation for uplink OFDM systems
- **Algorithm:** Primal-Dual algorithm
- **Properties**
 - ▶ **Distributed** and **simple** updates
 - ▶ **Low** communication overhead per iteration
- **Future Work:**
 - ▶ Convergence with **delay** and **asynchronous** updates
 - ▶ **Faster** convergence

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