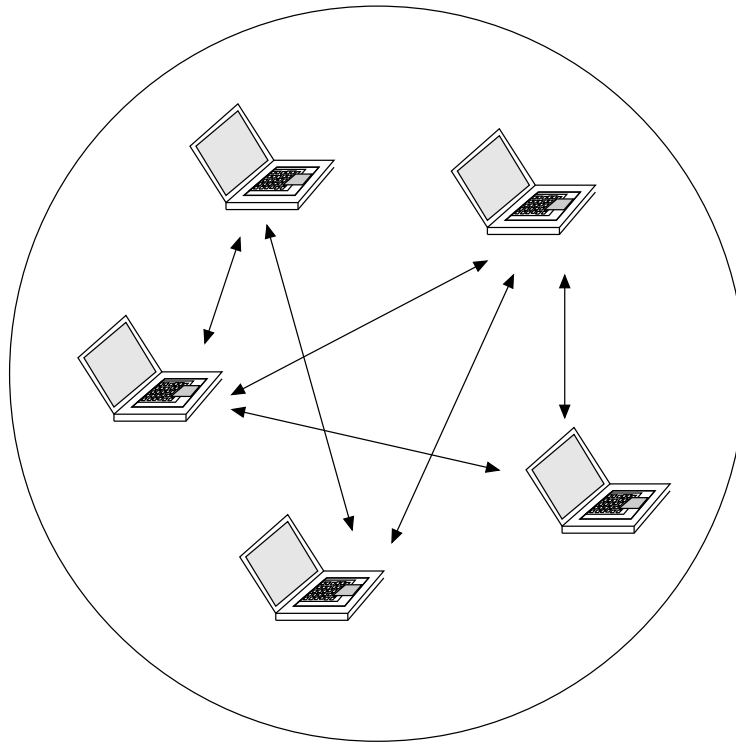


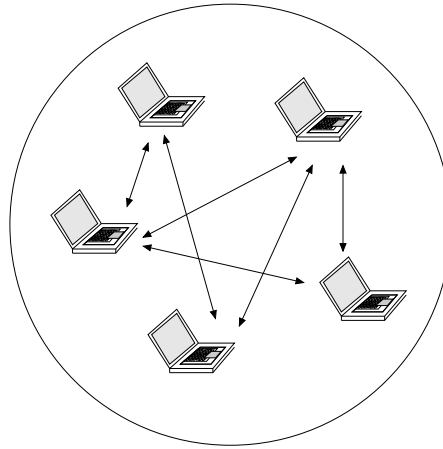
Active Queue Management and Scheduling in Wireless Networks: The Single Cell Case

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Single Cell Ad Hoc Networks



Rate and Medium Access Control



- Rate Control (TCP)
- Medium Access Control (IEEE 802.11)
- Objectives
 - Throughput
 - Fairness

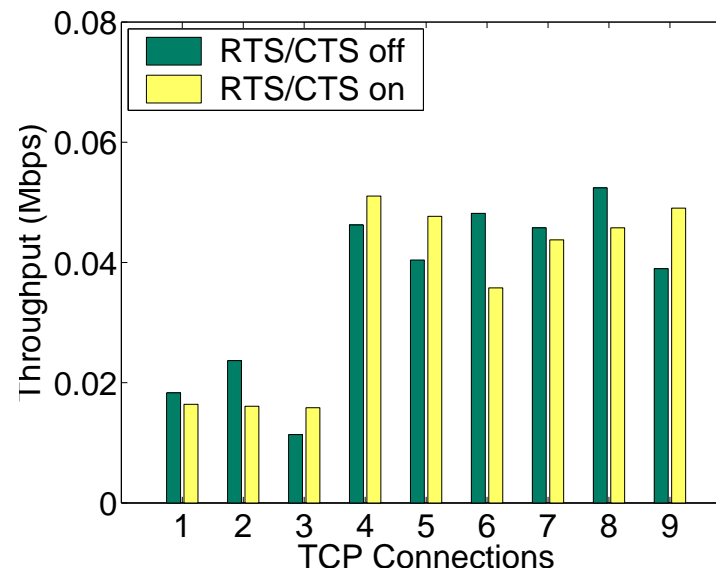
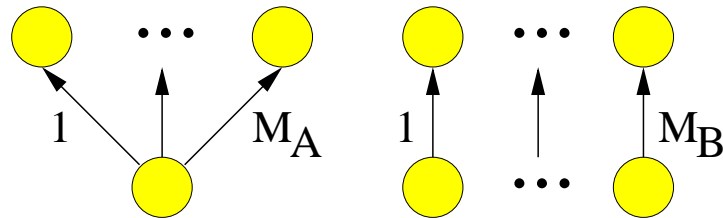
Wireline

- Rate Control: TCP – \rightarrow Stability
- AQM: RED, REM, BLUE – \rightarrow Throughput (Operating Point)
- Scheduling: FIFO – \rightarrow Fairness

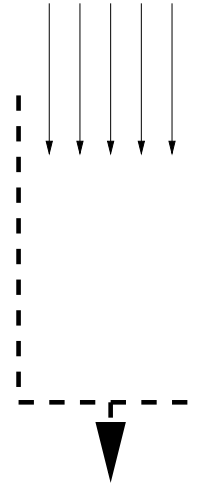
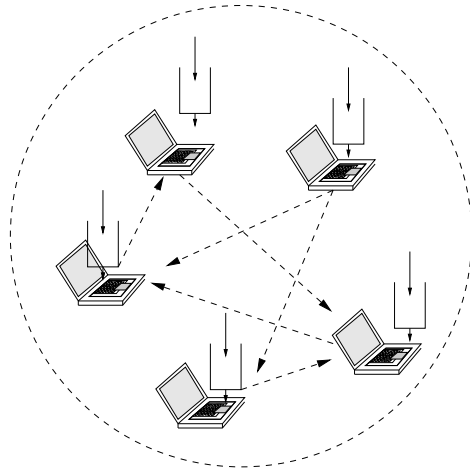
Outline

- Interaction of TCP Reno and IEEE 802.11
- Fair MAC Protocol
 - Properties
 - Centralized Solution
 - Distributed Solution
- Interaction with TCP
 - Single Cell Networks
- Discussion

TCP Reno and IEEE 802.11: Single Cell Network

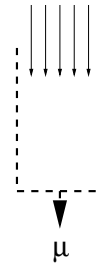
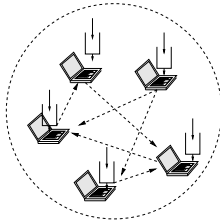


Properties of a Fair MAC Protocol



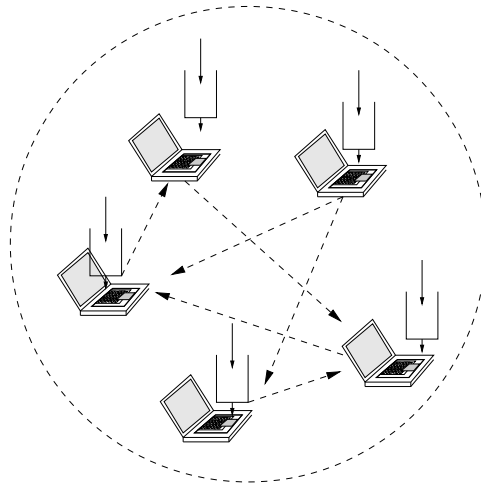
- All Packets
 - Same Delay
 - Same Drop Probability

Properties of a Fair MAC Protocol



- Network Arrival Rate $\lambda = \sum \lambda_n$
- Network Throughput $X(\lambda)$
- Distributed Buffer with Service Rate μ
 - All Packets have same Delay
 - All Packet experience same Drop Probability
 - $X(\lambda)$ increasing in λ
 - $\lim_{\lambda \rightarrow \infty} X(\lambda) = \mu$

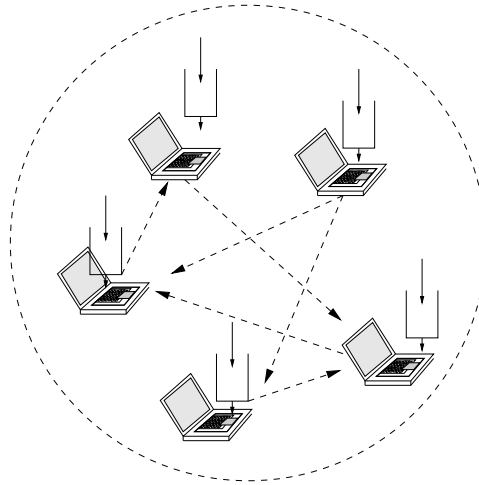
Centralized Solution



- Know Backlog b_n at each node n
- Total Backlog $B = \sum b_n$
- Schedule Node n with Probability

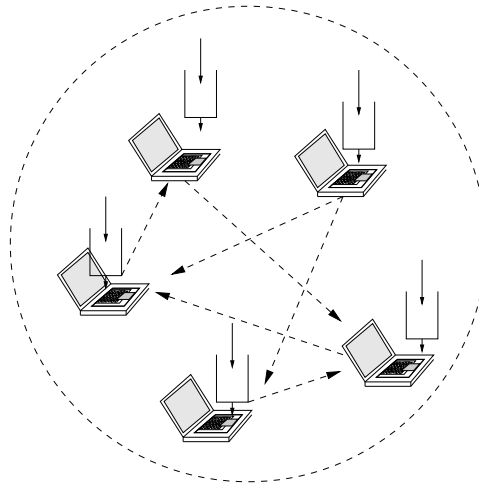
$$q_n = \frac{b_n}{B}$$

Centralized Solution



- Infinite Buffers
- Poisson Arrivals with Rate λ
- Exponential Transmission Delay with Mean $1/\mu$
- Performance as in $M/M/1$ queue

Distributed Solution

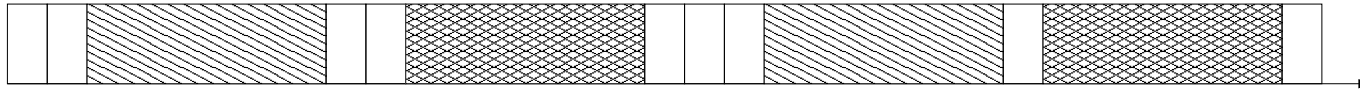


- CSMA
- Node n makes Transmission Attempt with Probability

$$q_n = \min\{b_n q, 1 - \epsilon\}$$

- CSMA with Backlog-Dependent Transmission Probabilities

Distributed Solution



- Idle Slot, L_i

$$P_i = \prod_{n=1}^N (1 - qb_n) \approx \prod_{n=1}^N e^{-qb_n} = e^{-qB} = e^{-G}$$

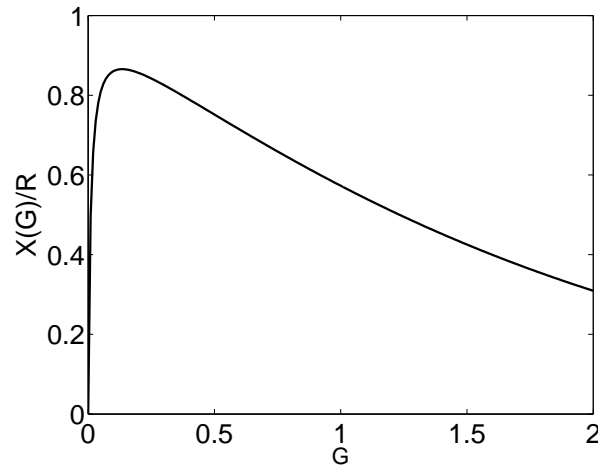
- Offered Load: $G = qB = \sum qb_n$
- Successful Transmission, L_p

$$P_s = Ge^{-G}$$

- Throughput

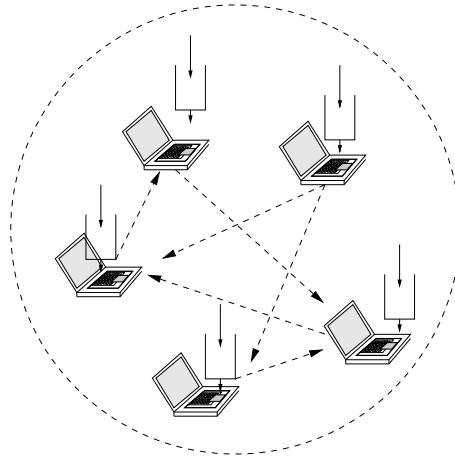
$$X(G) = \frac{Ge^{-G}}{L_i + (1 - e^{-G})L_p}$$

Distributed Solution



- Throughput $X(G) = \frac{Ge^{-G}}{L_i + (1 - e^{-G})L_p}$
- $G^+ = \sqrt{\frac{2L_i}{L_p}}$
- $G = qB = q \sum b_n$
- Active Queue Management: Keep Total Backlog at $B^* = G^*/q$

Active Queue Management



- Packet-Drop Probability $p(u)$

$$p(u) = \begin{cases} \kappa u, & 0 \leq u \leq 1/\kappa, \\ 1, & u > 1/\kappa, \end{cases}$$

- Use Channel Feedback to Compute u

Active Queue Management

- How to adapt u ?
 - Idle Slot: Decrease by α
 - Busy Slot: Increase by β
- Expected Change under $G = qB$

$$\Delta u = -\alpha P_i + (-\alpha + \beta)P_b = -\alpha + \beta(1 - e^{-G}),$$

- Operating Point G^* : $\Delta u = 0$

$$-\alpha + \beta(1 - e^{-G^*}) = 0$$

or

$$G^* = \ln \left(\frac{\beta}{\beta - \alpha} \right)$$

Performance

- Throughput $X(G^*)$:

$$\alpha = \beta(1 - e^{-G^*})$$

- Backlog B^*

$$q = \frac{G^*}{B^*}$$

- Implements Distributed Queue with Rate $X(G^*)$
- B^* can be Chosen

Summary

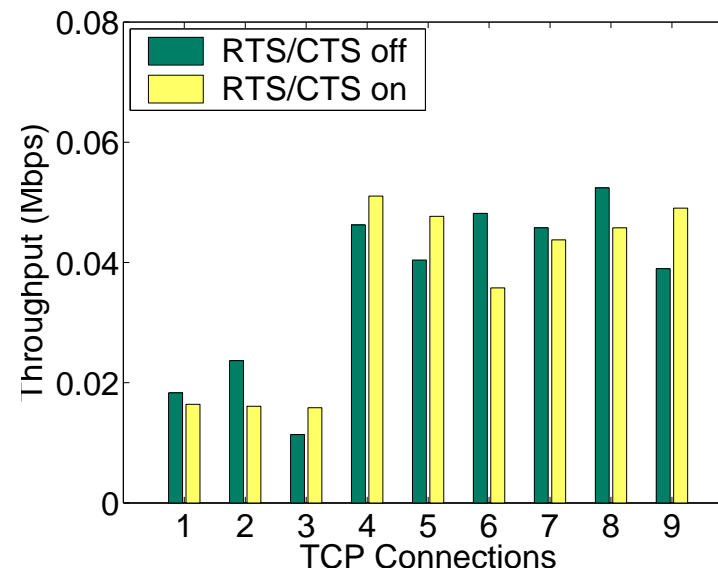
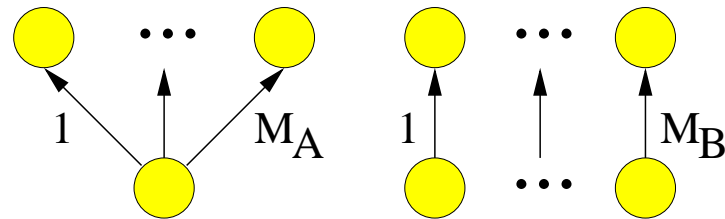
- Scheduling: CSMA with Backlog-Dependent Transmission Probabilities

$$q_n = qb_n$$

- Active Queue Management: Drop Probability $p(u)$
 - Idle Slot: Decrease u by α
 - Busy Slot: Increase u by β
- Design of Virtual Buffer
 - Throughput $X(G^*)$
 - Buffer Size B^*

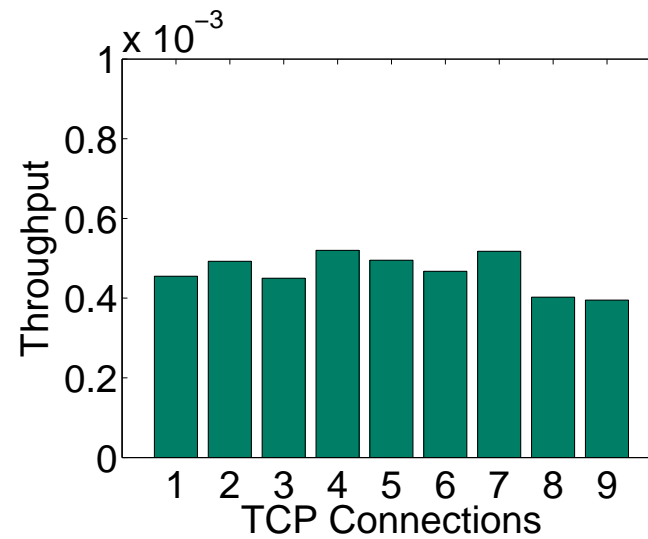
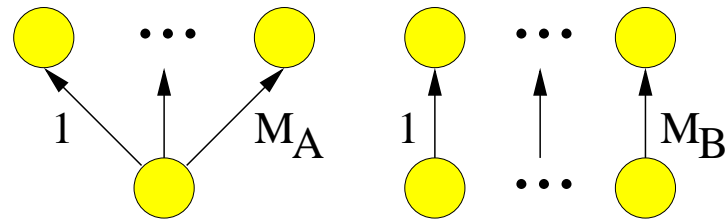
Interaction with TCP Reno: Single Cell Case

- TCP Reno and 802.11



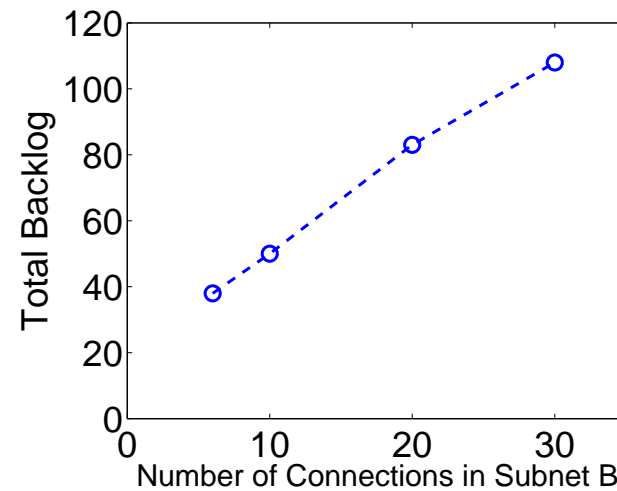
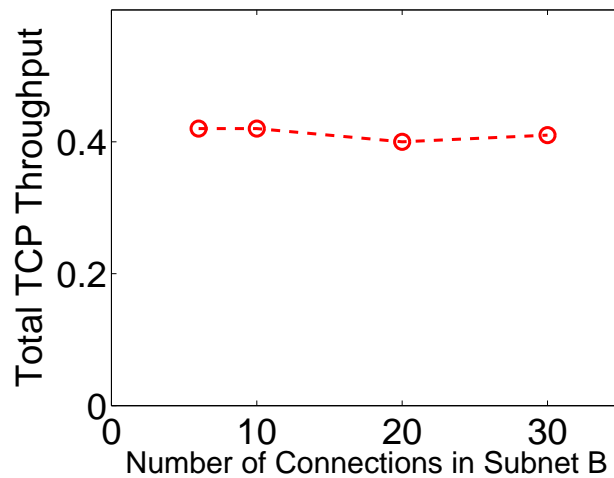
Interaction with TCP Reno: Single Cell Case

- TCP Reno and with New Mechanism



Interaction with TCP Reno: Single Cell

- Scaling as M_B increases
- $B^* = O(N)$
- $X(G^+) = 0.42$



Modelling Interaction with TCP Reno

- M Connections
- D_m : RTT of connection m
- Throughput

$$\begin{aligned} & \max \sum_{m=1}^M \frac{\sqrt{2}}{D_m} \arctan \left(\frac{x_m D_m^2}{2} \right) \\ & s.t. \sum_m x_m \leq X(G^*) \\ & \quad x_m \geq 0, m = 1, \dots, M, \end{aligned}$$

- $U_m(x_m) = \frac{\sqrt{2}}{D_m} \arctan \left(\frac{x_m D_m^2}{2} \right)$

Related Work

- AQM: RED, REM, BLUE
- NRED: RED in Wireless Networks with 802.11
- Utility Maximization
- Systematic - Practical

Conclusion

- Rate Control and MAC in Wireless Networks
- Protocol Design in Wireline Networks
- Layers
- IEEE 802.11
- Multihop Networks