

**This lecture:**

- Convex optimization
  - Convex sets
  - Convex functions
  - Convex optimization problems
  - Why convex optimization? Why so early in the course?

Recall the general form of our optimization problems:

$$\begin{array}{ll} \min. & f(x) \\ \text{s.t.} & x \in \Omega \end{array}$$

- In the last lecture, we focused on unconstrained optimization:  $\Omega = \mathbb{R}^n$ .
- We saw the definitions of local and global optimality, and, first and second order optimality conditions.
- In this lecture, we consider a very important special case of constrained optimization problems known as "*convex optimization problems*".
- For these problems,
  - $f$  will be a "*convex function*".
  - $\Omega$  will be a "*convex set*".
  - These notions are defined formally in this lecture.
- Roughly speaking, the high-level message is this:
  - Convex optimization problems are pretty much the broadest class of optimization problems that we know how to solve efficiently.
  - They have nice geometric properties;
    - e.g., a local minimum is automatically a global minimum.
  - Numerous important optimization problems in engineering, operations research, machine learning, etc. are convex.
  - There is available software that can take (a large subset of) convex problems written in very high-level language and solve it.
    - You should take advantage of this!
  - Convex optimization is one of the biggest success stories of modern theory of optimization.

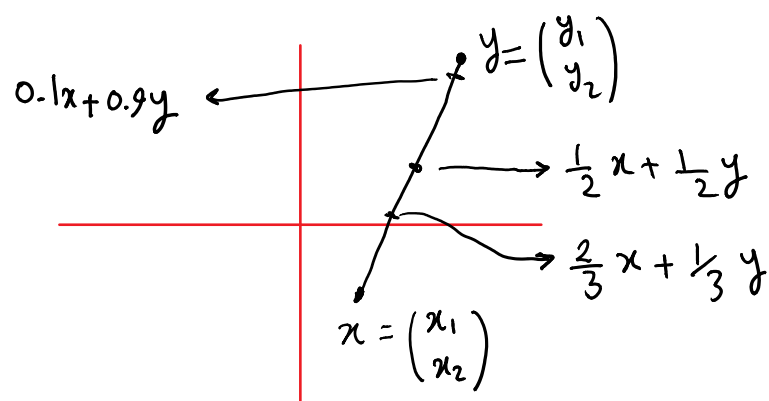
## Convex sets

**Definition.** A set  $\Omega \subseteq \mathbb{R}^n$  is *convex*, if for all  $x, y \in \Omega$ , the line segment connecting  $x$  and  $y$  is also in  $\Omega$ . In other words,

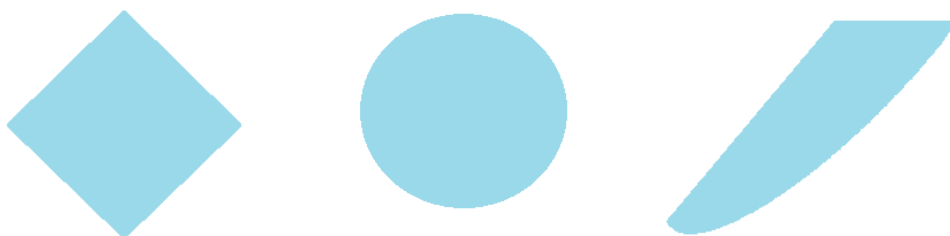
$$x, y \in \Omega, \lambda \in [0, 1] \Rightarrow \lambda x + (1 - \lambda)y \in \Omega$$

- A point of the form  $\lambda x + (1 - \lambda)y$ ,  $\lambda \in [0, 1]$  is called a *convex combination* of  $x$  and  $y$ .
- Note that when  $\lambda = 0$ , we are at  $y$ ; when  $\lambda = 1$ , we are at  $x$ ; for intermediate values of  $\lambda$ , we are on the line segment connecting  $x$  and  $y$ .

Illustration of the concept of a convex combination:



**Convex:**



**Not convex:**



## Convex sets & Midpoint Convexity

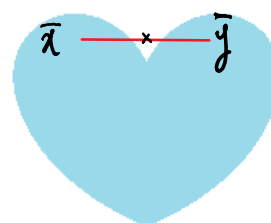
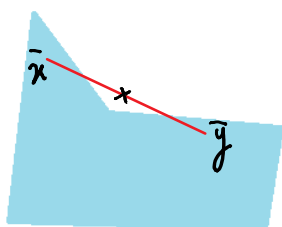
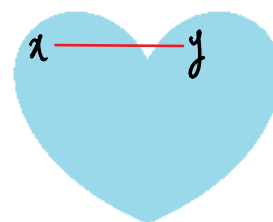
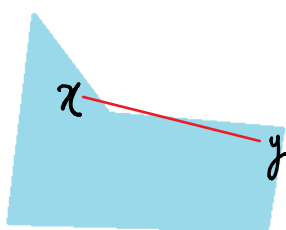
Midpoint convexity is a notion that is equivalent to convexity in most practical settings, but it is a little bit cleaner to work with.

**Definition.** A set  $\Omega \subseteq \mathbb{R}^n$  is *midpoint convex*, if for all  $x, y \in \Omega$ , the midpoint between  $x$  and  $y$  is also in  $\Omega$ . In other words,

$$x, y \in \Omega \Rightarrow \frac{1}{2}x + \frac{1}{2}y \in \Omega.$$

- Obviously, convex sets are midpoint convex.
- Under mild conditions, midpoint convex sets are convex
  - e.g., a closed midpoint convex sets is convex.
  - What is an example of a midpoint convex set that is not convex?  
(The set of all rational points in  $[0,1]$ .)

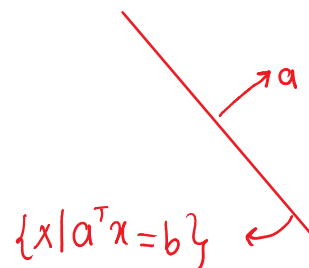
The nonconvex sets that we had are also not midpoint convex (why)?:



## Common convex sets in optimization

(Prove convexity in each case.)

- **Hyperplanes:**  $\{x \mid a^T x = b\}$  ( $a \in \mathbb{R}^n, b \in \mathbb{R}, a \neq 0$ )

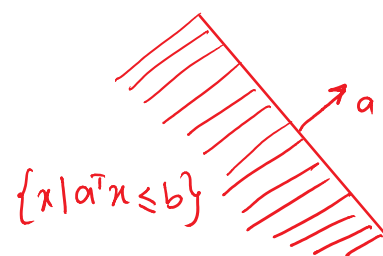


- **Halfspaces:**  $\{x \mid a^T x \leq b\}$  ( $a \in \mathbb{R}^n, b \in \mathbb{R}, a \neq 0$ )

Proof: Let  $\mathcal{H} := \{x \mid a^T x \leq b\}$ . Take  $x, y \in \mathcal{H}$

$$a^T (\lambda x + (1-\lambda)y) = \lambda a^T x + (1-\lambda)a^T y \leq \lambda b + (1-\lambda)b = b$$

$$\Rightarrow \lambda x + (1-\lambda)y \in \mathcal{H}. \quad \square$$



- **Euclidean balls:**  $\{x \mid \|x - x_c\| \leq r\}$  ( $x_c \in \mathbb{R}^n, r \in \mathbb{R}, \|\cdot\|$  2-norm)

Proof: Let  $\mathcal{B} := \{x \mid \|x - x_c\| \leq r\}$ . Take  $x, y \in \mathcal{B}$ .

$$\|\lambda x + (1-\lambda)y - x_c\| = \|\lambda(x - x_c) + (1-\lambda)(y - x_c)\|$$

$$\stackrel{\text{Triangle ineq.}}{\leq} \|\lambda(x - x_c)\| + \|(1-\lambda)(y - x_c)\| \stackrel{\text{homog.}}{=} \lambda\|x - x_c\| + (1-\lambda)\|y - x_c\| \stackrel{x, y \in \mathcal{B}}{\leq} \lambda r + (1-\lambda)r = r. \Rightarrow \lambda x + (1-\lambda)y \in \mathcal{B}. \quad \square$$



- **Ellipsoids:**  $\{x \mid (x - x_c)^T P (x - x_c) \leq r\}$  ( $x_c \in \mathbb{R}^n, r \in \mathbb{R}, P \succ 0$ )

( $P$  here is an  $n \times n$  symmetric matrix)

Proof hint: Wait until you see convex functions and quasiconvex functions. Observe that ellipsoids are sublevel sets of convex quadratic functions.



## Fancier convex sets

Many fundamental objects in mathematics have surprising convexity properties.

For example, prove that the following two sets are convex.

- The set of (symmetric) positive semidefinite matrices:  
 $S_+^{n \times n} = \{P \in S^{n \times n} \mid P \succcurlyeq 0\}$

*Proof.* Let  $A \succcurlyeq 0, B \succcurlyeq 0$ . Let  $\lambda \in [0, 1]$ .  $x^T (\lambda A + (1-\lambda)B)x = \underbrace{\lambda x^T A x}_{\geq 0} + \underbrace{(1-\lambda)x^T B x}_{\geq 0} \geq 0. \square$

e.g.,  $\{x, y, z \mid \begin{bmatrix} x & y \\ y & z \end{bmatrix} \succcurlyeq 0\}$ :

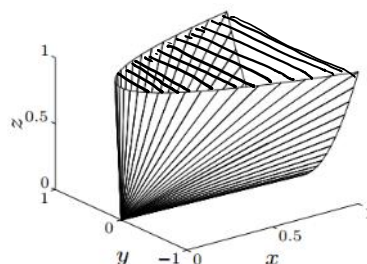
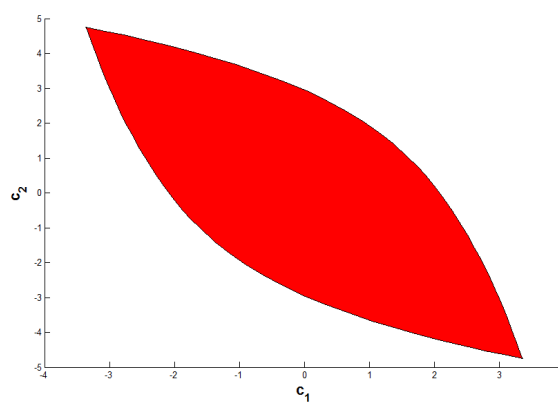


Image credit: [BV04]

- The set of nonnegative polynomials in  $n$  variables and of degree  $d$ .  
 (A polynomial  $p(x_1, \dots, x_n)$  is nonnegative, if  $p(x) \geq 0, \forall x \in \mathbb{R}^n$ .)

e.g.,  $\{(c_1, c_2) \mid 2x_1^4 + x_2^4 + c_1x_1x_2^3 + c_2x_1^3x_2 \geq 0, \forall (x_1, x_2) \in \mathbb{R}^2\}$ :



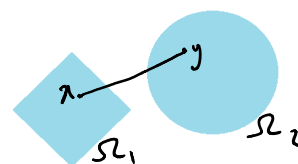
## Intersections of convex sets

- Easy to see that intersection of two convex sets is convex:  
 $\Omega_1$  convex,  $\Omega_2$  convex  $\Rightarrow \Omega_1 \cap \Omega_2$  convex.

Proof:

Pick  $x \in \Omega_1 \cap \Omega_2$ ,  $y \in \Omega_1 \cap \Omega_2$   
 $\forall \lambda \in [0,1]$ ,  $\lambda x + (1-\lambda)y \in \Omega_1$  (b/c  $\Omega_1$  is convex)  
 $\left\{ \begin{array}{l} \lambda x + (1-\lambda)y \in \Omega_2 \text{ (b/c } \Omega_2 \text{ is convex)} \\ \lambda x + (1-\lambda)y \in \Omega_1 \end{array} \right. \Rightarrow \lambda x + (1-\lambda)y \in \Omega_1 \cap \Omega_2 \quad \square$

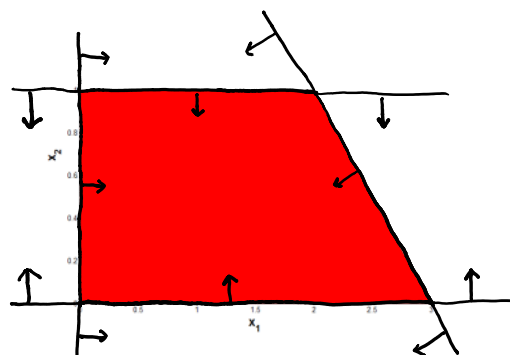
- Obviously, the union may not be convex:



## Polyhedra

- A polyhedron is the solution set of finitely many linear inequalities.
  - Ubiquitous in optimization theory.
  - Feasible sets of "linear programs" (an upcoming subject).
- Such sets are written in the form:
 
$$\{x \mid Ax \leq b\},$$
 where  $A$  is an  $m \times n$  matrix, and  $b$  is an  $m \times 1$  vector.
- These sets are convex: intersection of halfspaces  $a_i^T x \leq b_i$ ,  
 where  $a_i^T$  is the  $i$ -th row of  $A$ .

e.g.,  $A = \begin{bmatrix} -1 & 0 \\ 0 & -1 \\ 0 & 1 \\ 1 & 1 \end{bmatrix}, b = \begin{bmatrix} 0 \\ 0 \\ 1 \\ 3 \end{bmatrix}$

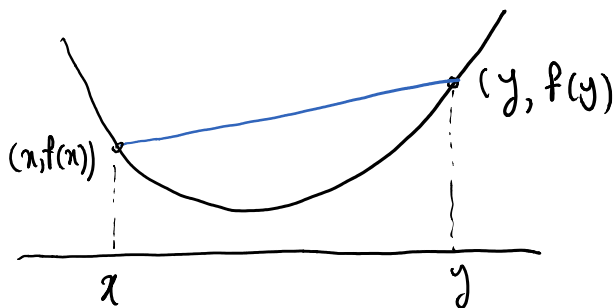


## Convex functions

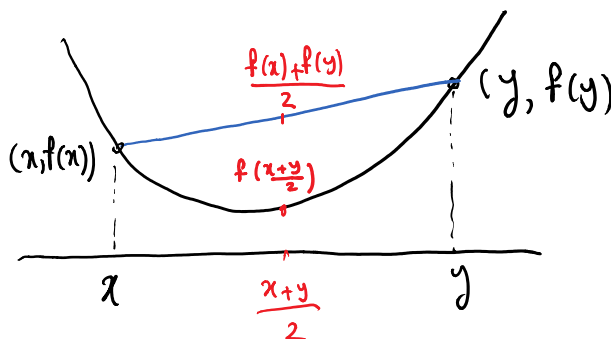
**Definition.** A function  $f: \mathbb{R}^n \rightarrow \mathbb{R}$  is *convex* if its domain is a convex set and for all  $x, y$  in its domain, and all  $\lambda \in [0, 1]$ , we have

$$f(\lambda x + (1 - \lambda)y) \leq \lambda f(x) + (1 - \lambda)f(y).$$

- In words: take any two points  $x, y$ ;  $f$  evaluated at any convex combination should be no larger than the same convex combination of  $f(x)$  and  $f(y)$ .
- If  $\lambda = \frac{1}{2}$ , interpretation is even easier: take any two points  $x, y$ ;  $f$  evaluated at the midpoint should be no larger than the average of  $f(x)$  and  $f(y)$ .
- Geometrically, the line segment connecting  $(x, f(x))$  to  $(y, f(y))$  sits above the graph of  $f$ .



$(f: \mathbb{R} \rightarrow \mathbb{R})$



**Definition.** A function  $f: \mathbb{R}^n \rightarrow \mathbb{R}$  is

- *Concave*, if  $\forall x, y, \forall \lambda \in [0,1]$   
 $f(\lambda x + (1 - \lambda)y) \geq \lambda f(x) + (1 - \lambda)f(y)$ .
- *Strictly convex*, if  $\forall x, y, x \neq y, \forall \lambda \in (0,1)$   
 $f(\lambda x + (1 - \lambda)y) < \lambda f(x) + (1 - \lambda)f(y)$ .
- *Strictly concave*, if  $\forall x, y, x \neq y, \forall \lambda \in (0,1)$   
 $f(\lambda x + (1 - \lambda)y) > \lambda f(x) + (1 - \lambda)f(y)$ .

**Note:**  $f$  is concave if and only if  $-f$  is convex. Similarly,  $f$  is strictly concave if and only if  $-f$  is strictly convex.

The only functions that are both convex and concave are affine functions; i.e., functions of the form:

$$f(x) = a^T x + b, \quad (a \in \mathbb{R}^n, b \in \mathbb{R}).$$



convex  
(and strictly convex)



concave  
(and strictly concave)



neither convex  
nor concave



both convex and  
concave (but not  
strictly)

Let's see some examples of convex functions (selection from [BV04]; see this reference for many more examples).

**Examples of univariate convex functions ( $f: \mathbb{R} \rightarrow \mathbb{R}$ ):**

- $e^{ax}$
- $-\log x$
- $x^a$  (defined on  $\mathbb{R}_{++}$ )  $a \geq 1$  or  $a \leq 0$
- $-x^a$  (defined on  $\mathbb{R}_{++}$ )  $0 \leq a \leq 1$
- $|x|^a, a \geq 1$
- $x \log x$  (defined on  $\mathbb{R}_{++}$ )
- Try to plot the functions above and convince yourself of convexity visually.
- Can you formally verify that these functions are convex?
- We will soon see some characterizations of convex functions that make the task of verifying convexity a bit easier.



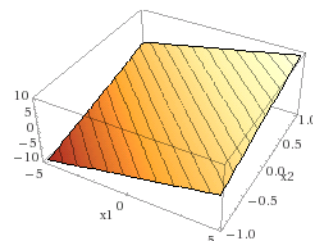
## Examples of convex functions ( $f: \mathbb{R}^n \rightarrow \mathbb{R}$ )

- **Affine functions:**  $f(x) = a^T x + b$  (for any  $a \in \mathbb{R}^n, b \in \mathbb{R}$ )

(convex, but not strictly convex; also concave)

$$\text{Proof: } \forall \lambda \in [0, 1], f(\lambda x + (1-\lambda)y) = a^T(\lambda x + (1-\lambda)y) + b$$

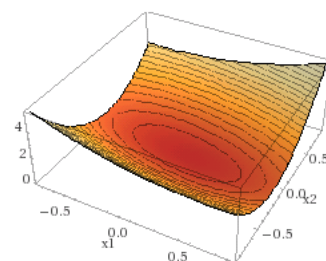
$$= \lambda a^T x + (1-\lambda) a^T y + \lambda b + (1-\lambda)b = \lambda f(x) + (1-\lambda) f(y). \quad \square$$



- **Some quadratic functions:**

$$f(x) = x^T Q x + c^T x + d$$

- Convex if and only if  $Q \succeq 0$ .
- Strictly convex if and only if  $Q \succ 0$ .
- Concave iff  $Q \preceq 0$ ; Strictly concave iff  $Q \prec 0$ .
- Proofs are easy from the second order characterization of convexity (coming up).



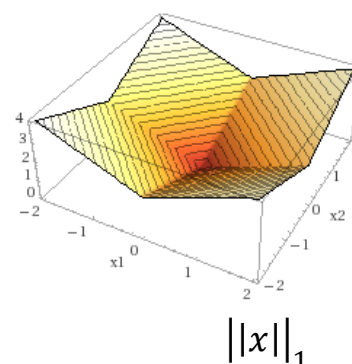
- **Any norm:** meaning, any function  $f$  satisfying:

- $f(\alpha x) = |\alpha| f(x), \forall \alpha \in \mathbb{R}$
- $f(x + y) \leq f(x) + f(y)$
- $f(x) \geq 0, \forall x, f(x) = 0 \Rightarrow x = 0$

$$\forall \lambda \in [0, 1]$$

$$\text{Proof: } f(\lambda x + (1-\lambda)y) \stackrel{b}{\leq} f(\lambda x) + f((1-\lambda)y)$$

$$\stackrel{a}{=} \lambda f(x) + (1-\lambda) f(y). \quad \square$$



Examples:

- $\|x\|_\infty = \max_i |x_i|$
- $\|x\|_p = \left( \sum_{i=1}^n |x_i|^p \right)^{\frac{1}{p}}, \quad p \geq 1$
- $\|x\|_Q = \sqrt{x^T Q x}, Q \succ 0$

## Midpoint convex functions

Same idea as what we saw for midpoint convex sets.

**Definition.** A function  $f: \mathbb{R}^n \rightarrow \mathbb{R}$  is *midpoint convex* if its domain is a convex set and for all  $x, y$  in its domain, we have

$$f\left(\frac{x+y}{2}\right) \leq \frac{1}{2}f(x) + \frac{1}{2}f(y).$$

- Obviously, convex functions are midpoint convex.
- Continuous, midpoint convex functions are convex.

## Convexity = Convexity along all lines

**Theorem.** A function  $f: \mathbb{R}^n \rightarrow \mathbb{R}$  is convex if and only if the function  $g: \mathbb{R} \rightarrow \mathbb{R}$ , given by  $g(t) = f(x + ty)$  is convex (as a univariate function), for all  $x$  in domain of  $f$  and all  $y \in \mathbb{R}^n$ . (The domain of  $g$  here is all  $t$  for which  $x + ty$  is in the domain of  $f$ .)

- This should be intuitive geometrically:
  - The notion of convexity is defined based on line segments.
- The theorem simplifies many basic proofs in convex analysis.
- But it does not usually make verification of convexity that much easier; the condition needs to hold for *all* lines (and we have infinitely many).
- Many of the algorithms we will see in future lectures work by iteratively minimizing a function over lines. It's useful to remember that the restriction of a convex function to a line remains convex. Here is a proof:

Suppose for some  $x, y$ ,  $g(\alpha) = f(x + \alpha y)$  was not convex.

$$\Rightarrow \exists \lambda \in [0, 1], \alpha_1, \alpha_2 \text{ s.t. } g(\lambda \alpha_1 + (1-\lambda)\alpha_2) > \lambda g(\alpha_1) + (1-\lambda)g(\alpha_2).$$

$$\Rightarrow f\left(x + (\lambda \alpha_1 + (1-\lambda)\alpha_2)y\right) = f\left(\lambda(x + \alpha_1 y) + (1-\lambda)(x + \alpha_2 y)\right) > \lambda f(x + \alpha_1 y) + (1-\lambda)f(x + \alpha_2 y).$$

□

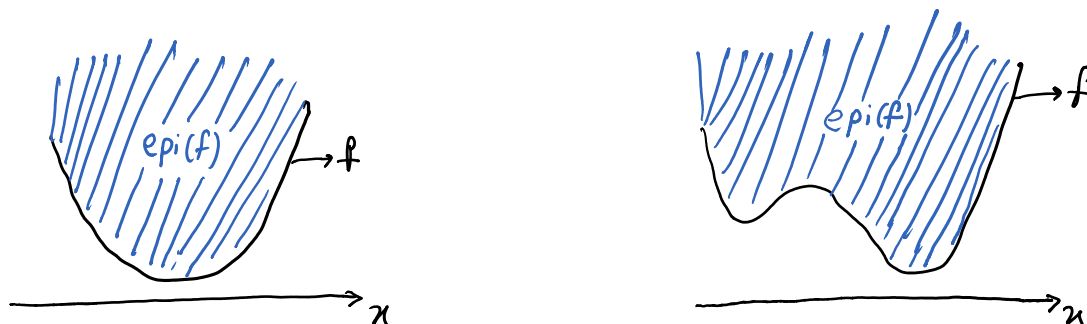
## Epigraph

Is there a connection between convex sets and convex functions?

- We will see a couple; via epigraphs, and sublevel sets.

**Definition.** The epigraph  $\text{epi}(f)$  of a function  $f: \mathbb{R}^n \rightarrow \mathbb{R}$  is a subset of  $\mathbb{R}^{n+1}$  defined as

$$\text{epi}(f) = \{(x, t) \mid x \in \text{domain}(f), f(x) \leq t\}.$$



**Theorem.** A function  $f: \mathbb{R}^n \rightarrow \mathbb{R}$  is convex if and only if its epigraph is convex (as a set).

*Proof:* Suppose  $f$  not convex  $\Rightarrow \exists x, y \in \text{dom}(f), \lambda \in [0, 1]$

$$\text{s.t. } f(\lambda x + (1-\lambda)y) > \lambda f(x) + (1-\lambda)f(y). \quad ①$$

Pick  $(x, f(x)), (y, f(y)) \in \text{epi}(f)$ .

$$① \Rightarrow (\lambda x + (1-\lambda)y, \lambda f(x) + (1-\lambda)f(y)) \notin \text{epi}(f).$$

Suppose  $\text{epi}(f)$  not convex  $\Rightarrow \exists (x, t_x), (y, t_y), \lambda \in [0, 1]$

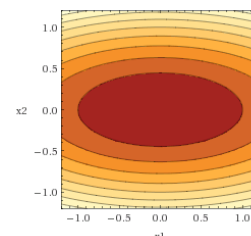
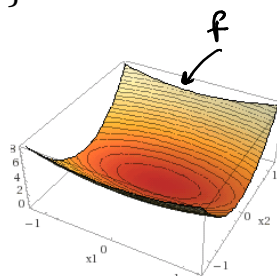
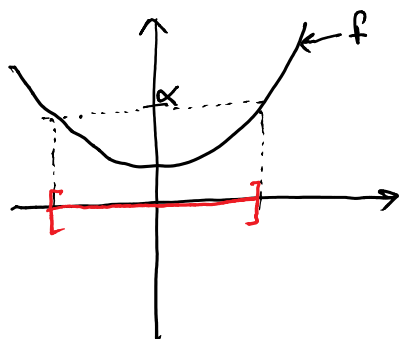
$$\text{s.t. } f(x) \leq t_x, \quad f(y) \leq t_y, \quad f(\lambda x + (1-\lambda)y) > \lambda t_x + (1-\lambda)t_y \\ > \lambda f(x) + (1-\lambda)f(y)$$

$\Rightarrow f$  not convex.  $\square$

## Convexity of sublevel sets

**Definition.** The  $\alpha$ -sublevel set of a function  $f: \mathbb{R}^n \rightarrow \mathbb{R}$  is the set

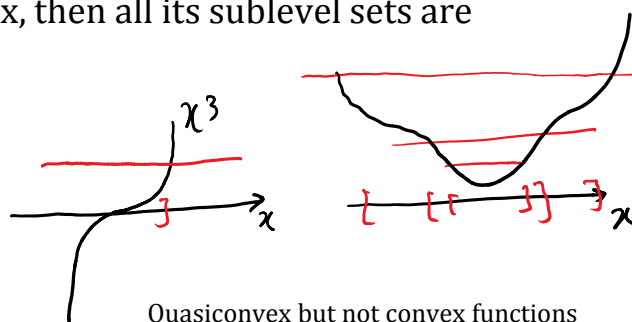
$$S_\alpha = \{x \in \text{domain}(f) \mid f(x) \leq \alpha\}.$$



Several sublevel sets (for different values of  $\alpha$ )

**Theorem.** If a function  $f: \mathbb{R}^n \rightarrow \mathbb{R}$  is convex, then all its sublevel sets are convex sets.

- Converse *not* true.
- A function whose sublevel sets are all convex is called *quasiconvex*.



Quasiconvex but not convex functions

Proof of theorem:

Pick  $x, y \in S_\alpha, \lambda \in [0, 1]$

$$x \in S_\alpha \Rightarrow f(x) \leq \alpha ; y \in S_\alpha \Rightarrow f(y) \leq \alpha$$

$$\begin{aligned} f \text{ convex} &\Rightarrow f(\lambda x + (1-\lambda)y) \leq \lambda f(x) + (1-\lambda)f(y) \\ &\leq \lambda \alpha + (1-\lambda)\alpha \\ &= \alpha \end{aligned}$$

$$\Rightarrow \lambda x + (1-\lambda)y \in S_\alpha. \quad \square$$

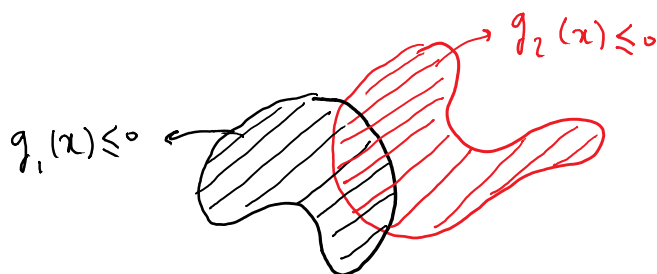
## Convex optimization problems

A convex optimization problem is an optimization problem of the form

$$\begin{aligned} \min. & f(x) \\ \text{s.t.} & g_i(x) \leq 0, \quad i = 1, \dots, m, \\ & h_j(x) = 0, \quad j = 1, \dots, k, \end{aligned}$$

where  $f, g_i: \mathbb{R}^n \rightarrow \mathbb{R}$  are convex functions and  $h_i: \mathbb{R}^n \rightarrow \mathbb{R}$  are affine functions.

- Let  $\Omega$  denote the feasible set:  $\Omega = \{x \in \mathbb{R}^n \mid g_i(x) \leq 0, h_j(x) = 0\}$ .
  - Observe that for a convex optimization problem  $\Omega$  is a convex set (why?)
  - But the converse is not true:
    - Consider for example,  $\Omega = \{x \in \mathbb{R} \mid x^3 \leq 0\}$ . Then  $\Omega$  is a convex set, but minimizing a convex function over  $\Omega$  is not a convex optimization problem per our definition.
    - However, the same set can be represented as  $\Omega = \{x \in \mathbb{R} \mid x \leq 0\}$ , and then this would be a convex optimization problem with our definition.
- Here is another example of a convex feasible set that fails our definition of a convex optimization problem:



$$\Omega = \{x \mid g_1(x) \leq 0, g_2(x) \leq 0\}$$

is a convex set. But neither

$g_1$  nor  $g_2$  are convex functions (why?).

## Convex optimization problems (cont'd)

- We require this stronger definition because otherwise many abstract and complex optimization problems can be formulated as optimization problems over a convex set. (Think, e.g., of the set of nonnegative polynomials.) The stronger definition is much closer to what we can actually solve efficiently.
- The software CVX that we'll be using ONLY accepts convex optimization problems defined as above.
- Beware that [CZ13] uses the weaker and more abstract definition for a convex optimization problem (i.e., the definition that simply asks  $\Omega$  to be a convex set.)

### Acceptable constraints in CVX:

- **Convex  $\leq$  Concave**
- **Affine  $==$  Affine**

This is really the same as:

- **Convex  $\leq 0$**
- **Affine  $== 0$**

Why?

(Hint: Sum of two convex functions is convex, and sums and differences of affine functions are affine. )

### Notes:

- Further reading for this lecture can include the first few pages of Chapters 2,3,4 of [BV04]. Your [CZ13] book defines convex sets in Section 4.3. Convex optimization appears in Chapter 22. The relevant sections are 22.1-22.3.

### References:

- [BV04] S. Boyd and L. Vandenberghe. Convex Optimization. Cambridge University Press, 2004.  
<http://stanford.edu/~boyd/cvxbook/>
- [CZ13] E.K.P. Chong and S.H. Zak. An Introduction to Optimization. Fourth edition. Wiley, 2013.