# ORF 363/COS 323 Computing and Optimization in the Physical and Social Sciences

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Princeton, ORFE

Lecture 1



### What is optimization?

- ■Roughly, can think of optimization as the science of making the most out of every situation.
- You've probably all done it many times recently:
- •What courses to take?
  - •To maximize learning.
  - ■To maximize GPA (?!)
  - Courses can't conflict.
  - ■Not before 10AM.
  - ■Professor rating > 4.5.

- •What furniture to buy?
  - ■To minimize cost.
  - To maximize comfort.
  - •Must fit in your room.
  - •Must have 3 drawers.
  - Not too heavy.

- •What path to choose for a run?
  - •Maximize encounter with other runners.
  - ■Path must be between 5 and 7 miles.
  - Chosen roads must have side walks.
  - ■Path should get you back home.

### **Common theme:**

- ■You make decisions and choose one of many alternatives.
- ■You hope to maximize or minimize something (you have an objective).
- ■You cannot make arbitrary decisions. Life puts constraints on you.



### How is this class different from your every-day optimization?

- •We'll be learning techniques for dealing with problems that have
  - ■Thousands (if not millions) of variables
  - ■Thousands (if not millions) of constraints
- ■These problems appear every day in the industry, in science, in engineering
- •Hopeless to make decisions in your head and with rules of thumb
- Need mathematical techniques that translate to algorithms
  - •Algorithms then get implemented on a computer to solve your optimization problem
- We typically model a physical or social scenario with a precise mathematical description
- ■In this mathematical model, we care about finding the best solution
- ■Whenever we can't find the best solution, we would like to know how far off our proposed solution is



### **Examples of optimization problems**

### In finance

- ■In what proportions to invest in 500 stocks?
  - ■To maximize return.
  - ■To minimize risk.
  - ■No more than 1/5 of your money in any one stock.
  - ■Transactions costs < \$70.
  - ■Return rate > 2%.

### In control engineering

- How to drive an autonomous vehicle from A to B?
  - ■To minimize fuel consumption.
  - ■To minimize travel time.
  - ■Distance to closest obstacle > 2 meters.
  - ■Speed < 40 miles/hr.
  - ■Path needs to be smooth (no sudden changes in direction).

### In machine learning

- •How to assign likelihoods to transactions being fraudulent?
  - ■To minimize probability of a false negative.
  - ■To penalize overfitting on training set.
  - ■Probability of false positive < .15.
  - Misclassification error on training set < 5%.</li>



### **Computing and Optimization**

- ■This class will give you a broad introduction to "optimization from a computational viewpoint."
- Optimization and computing are very close areas of applied mathematics:
  - ■For a host of major problems in computer science, the best algorithms currently come from the theory of optimization.
  - ■Conversely, foundational work by computer scientists has led to a shift of focus in optimization theory from "mathematical analysis" to "computational mathematics."
- ■Several basic topics in scientific computing (that we'll cover in this course) are either special cases or fundamental ingredients of more elaborate optimization algorithms:
  - ■Least squares, root finding, solving linear systems, solving linear inequalities, approximation and fitting, matrix factorizations, conjugate gradients,...



### Agenda for today

- Meet your teaching staff
- Get your hands dirty with algorithms
  - ■Game 1
  - ■Game 2
- Modelling with a mathematical program
  - Fermat's last theorem!
- Course logistics and expectations



# Meet your teaching staff (1/2)



**Amir Ali Ahmadi** (Amir Ali, or Amirali, is my first name)

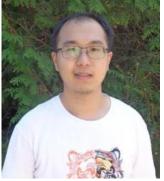
http://aaa.princeton.edu/ aaa@p...

- ■I am a Professor at ORFE, and affiliated faculty at COS, ECE, MAE, PACM, CSML
- ■I came to Princeton from MIT, EECS, after a fellowship at IBM Research
- I was on leave for a while, but now back to teaching!
- ■Office hours: Wed 4-6pm, Sherrerd 125



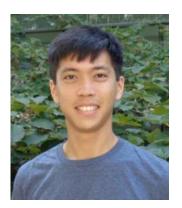
- Abraar Chaudhry (Honorary AI) ORFE senior grad student
- Has been an AI for this course several times in the past
- ■May teach a lecture or two
- Office hours: Wed 6-8pm, Sherrerd 122
- azc@p...

Meet your teaching staff (2/2)



■Yixuan Hua (AI) — ORFE grad student

- ■Office hours: Mon 4-6pm, Sherrerd 122
- ■yh7422@p...



■Jackie Lok (AI) – ORFE grad student

- ■Office hours: Tue 6-8pm, Sherrerd 122
- •jackie.lok@p...



■Pierfrancesco Beneventano (AI) – ORFE grad student

- ■Office hours: Mon 6-8pm, Sherrerd 122
- **■**pb29@p...



■Silu Li (AI) — ORFE grad student

- Office hours: Thu 6-8pm, Sherrerd 122
- ■sl3965@p...





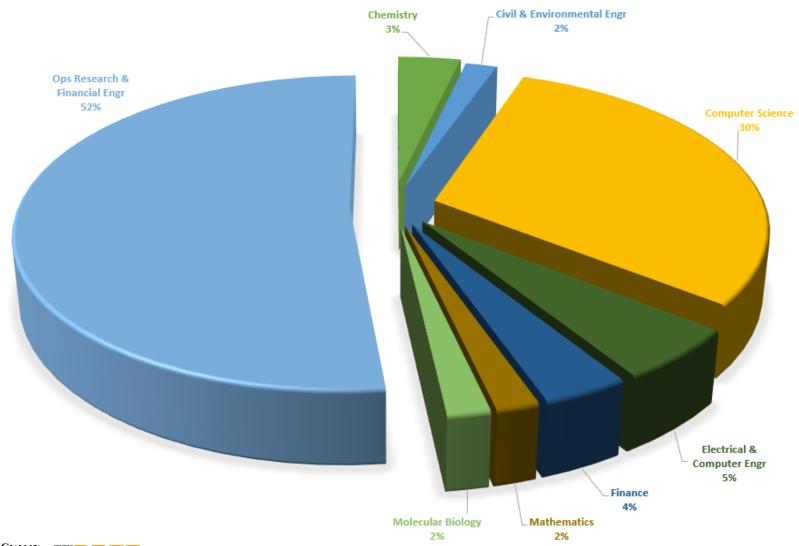
## **Office hours**

Mon	Tue	Wed	Thu	Fri
4-6pm Sherrerd 122 (Yixuan)		4-6pm Sherrerd 125 (AAA)		10am-12pm Sherrerd 122 (Als to rotate)
6-8pm Sherrerd 122 (Pier)	6-8pm Sherrerd 122 (Jackie)	6-8pm Sherrerd 122 (Abraar)	6-8pm Sherrerd 122 (Silu)	



# **Meet your classmates**

ORF 363/COS 323, Fall 2023 (56 registered students)





### **Course website**

### **Course website:**

http://aaa.princeton.edu/orf363

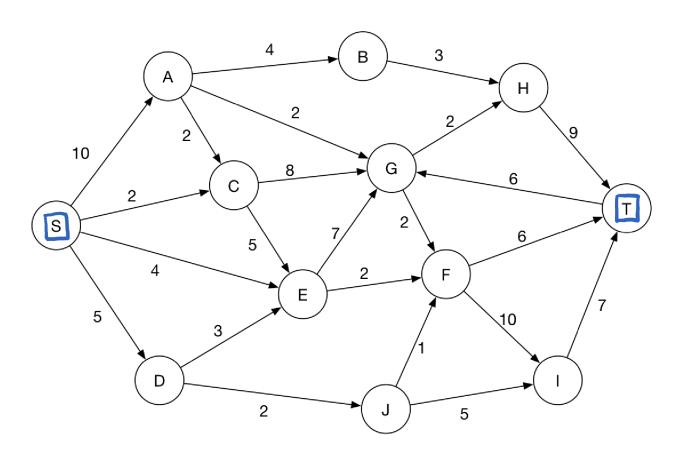
- Will have all lectures, problem sets, exams (a copy will also be posted on Canvas)
- Solutions will be posted only on Canvas
- Please set up notification settings on Canvas so you receive our announcements via email
- Please sign up to Ed Discussion via Canvas (and use it frequently!)
- Use Ed Discussion instead of email (unless your question is personal)



# Let's get to the games!



## Let's ship some oil together!

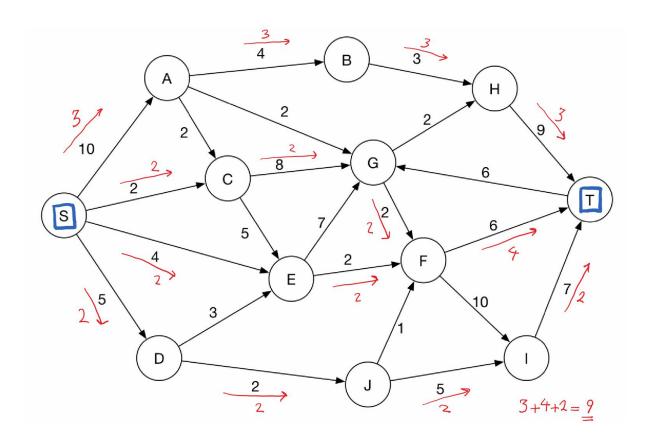


### Rules of the game:

- Cannot exceed capacity on the edges.
- ■For each node, except for S ant T, flow in = flow out (i.e., no storage).
- ■Goal: ship as much oil as you can from S to T.



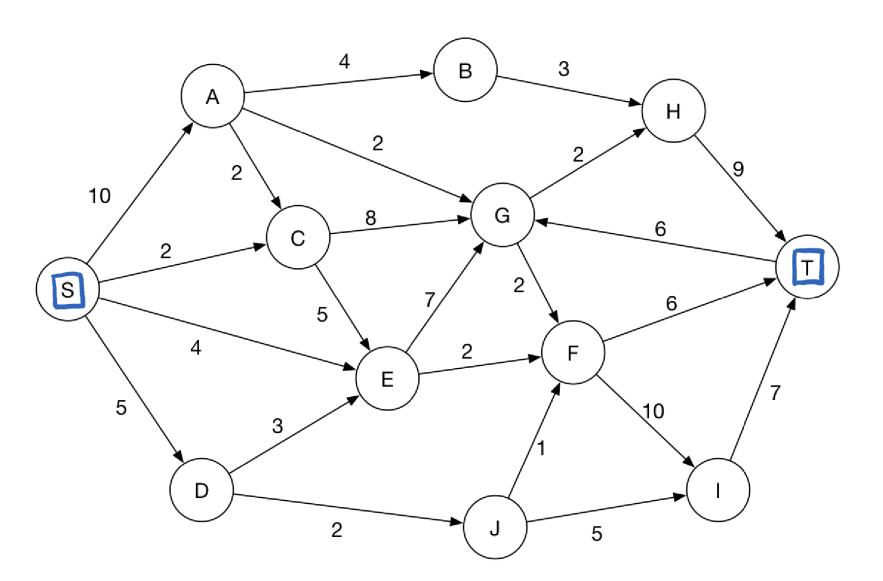
■Let me start things off for you. Here is a flow with value 9:



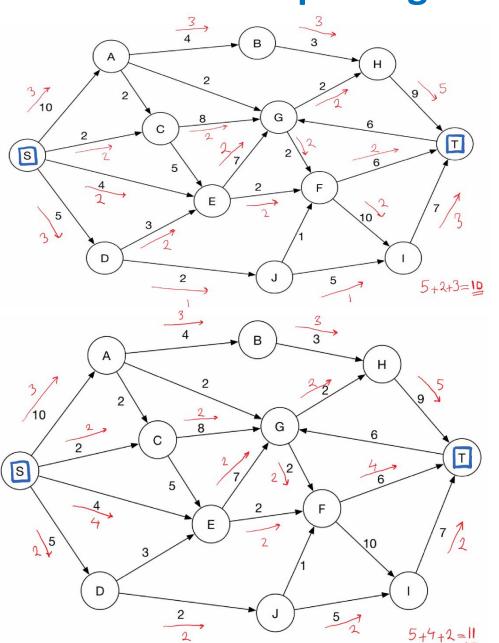
- ■Can you do better? How much better?
- ■You all get a copy of this graph.



# You tell me, I draw...







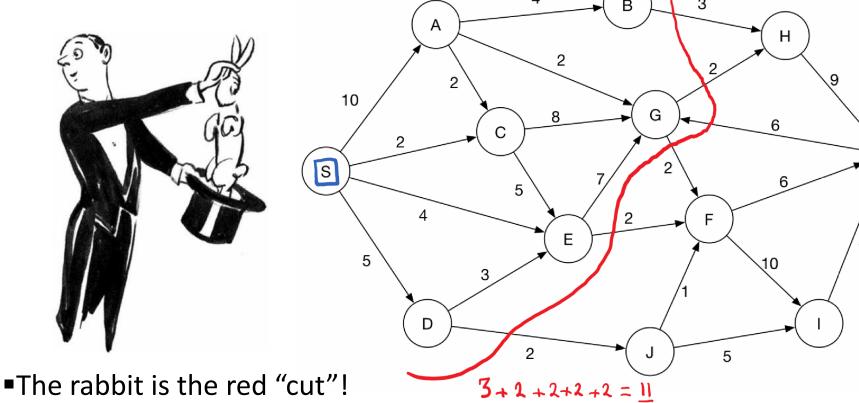
- ■Flow of value 10
- ■Can you do better?

- ■Flow of value 11
- Can you do better?
- •How can you prove that it's impossible to do better?



### 11 is the best possible!

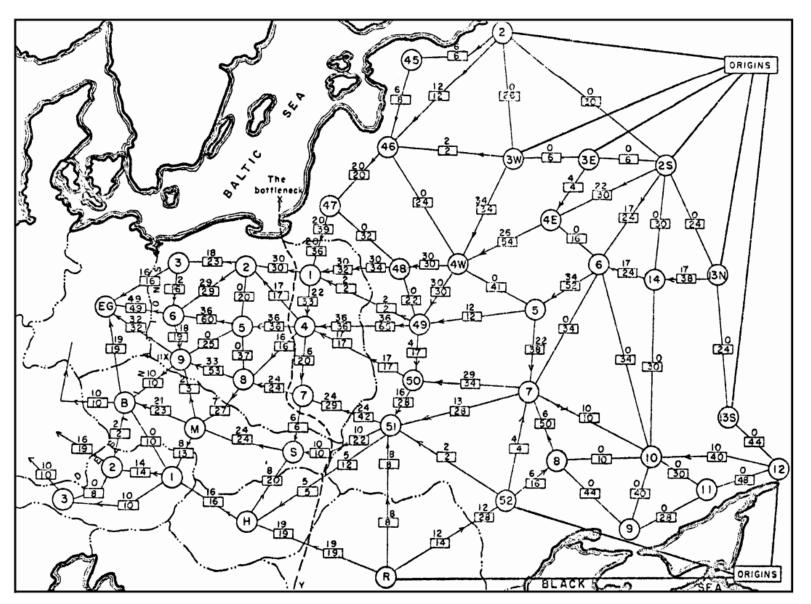
Proof by magic:



- Any flow from S to T must cross the red curve.
- ■So it can have value at most 11.



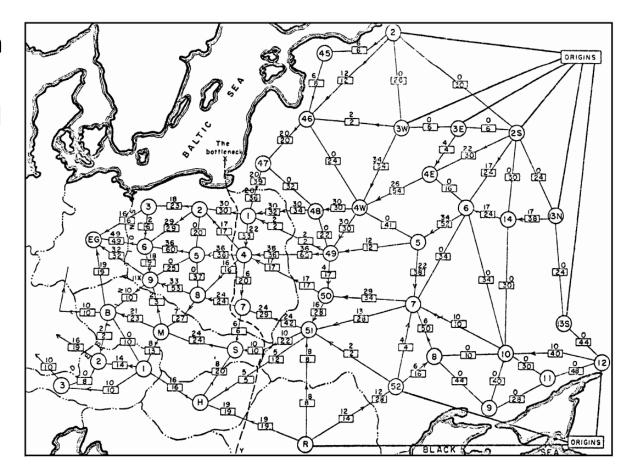
# Let's try a more realistic graph







- •How long do you think an optimization solver would take (on my laptop) to find the best solution here?
- •How many lines of code do you think you have to write for it?
- •How would someone who hasn't seen optimization approach this?

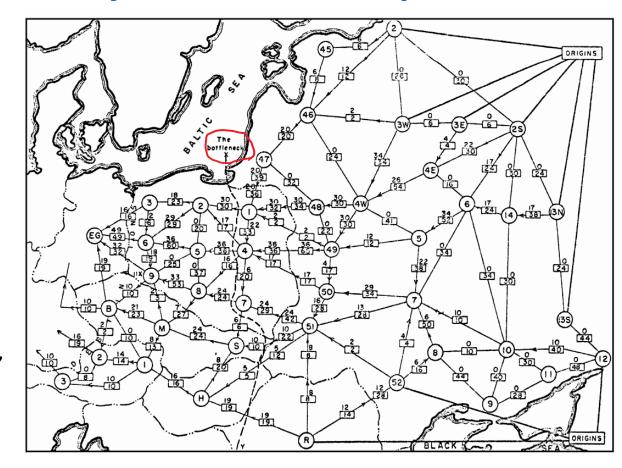


- Trial and error?
- ■Push a little flow here, a little there...
- Do you think they are likely to find the best solution?



### A bit of history behind this map

- ■From a secret report by Harris and Ross (1955) written for the Air Force.
- Railway network of the Western Soviet Union going to Eastern Europe.
- Declassified in 1999.
- ■Look at the min-cut on the map (called the "bottleneck")!
- ■There are 44 vertices, 105 edges, and the max flow is 163K.



- ■Harris and Ross gave a heuristic which happened to solve the problem optimally in this case.
- ■Later that year (1955), the famous Ford-Fulkerson algorithm came out of the RAND corporation. The algorithm always finds the best solution (for rational edge costs).



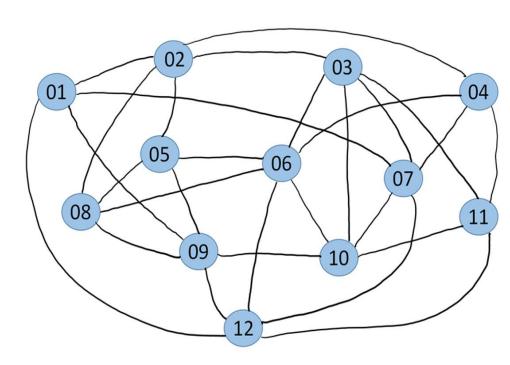
# Let's look at a second problem

...and tell me which one you thought was easier



## Two finals in one day? No thanks.

- ■The department chair at ORFE would like to schedule the final exams for 12 graduate courses offered this semester.
- ■He wants to have as many exams as possible on the same day, so everyone gets done quickly and goes on vacation.
- ■There is just one constraint:
  No student should have >1 exam on that day.
  - ■The nodes of this graph are the 12 courses.
  - ■There is an edge between two nodes if and only if there is at least one student who is taking both courses.
  - ■If we want to schedule as many exams as possible on the same day, what are we looking for in this graph?

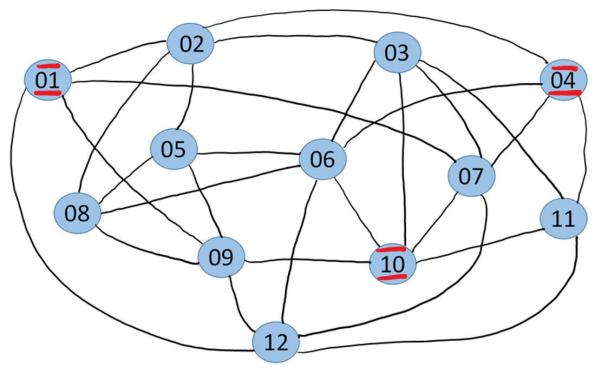


■The largest collection of nodes such that no two nodes share an edge.



Let me start things off for you. Here is 3 concurrent final

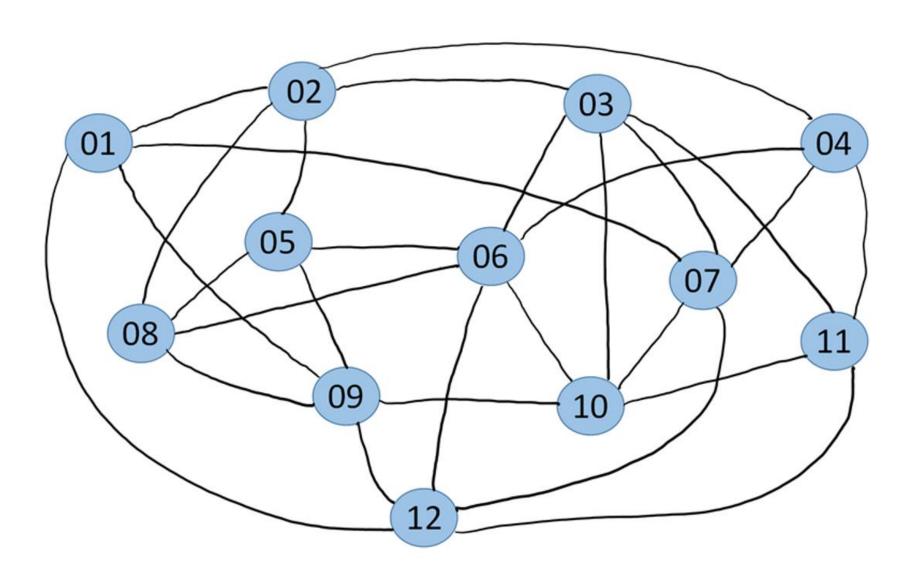
exams:



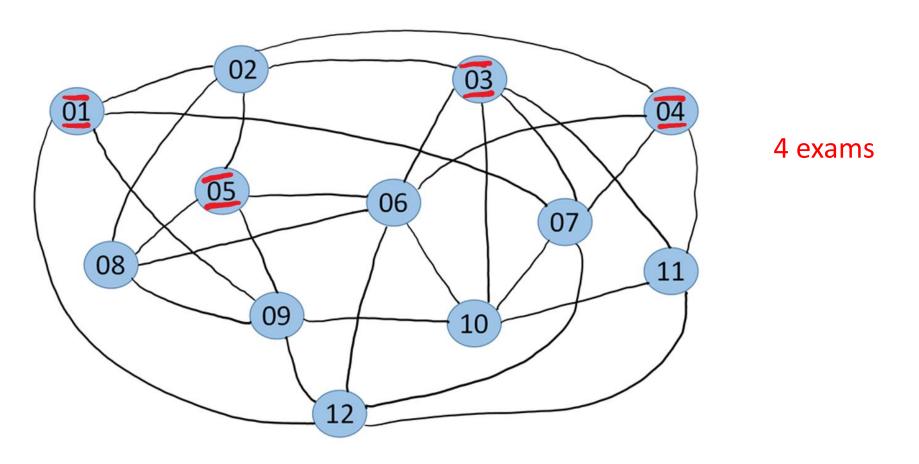
- Can you do better?
- ■How much better?
- ■You all get a copy of this graph.



# You tell me, I draw...

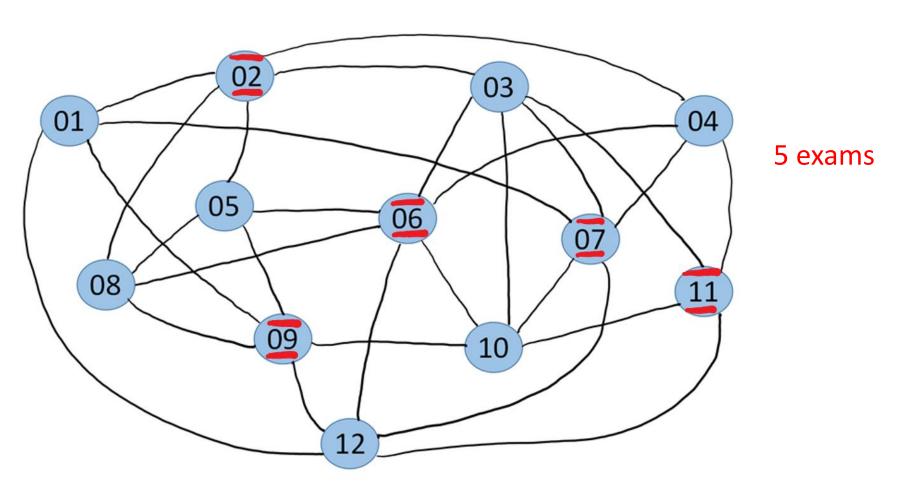






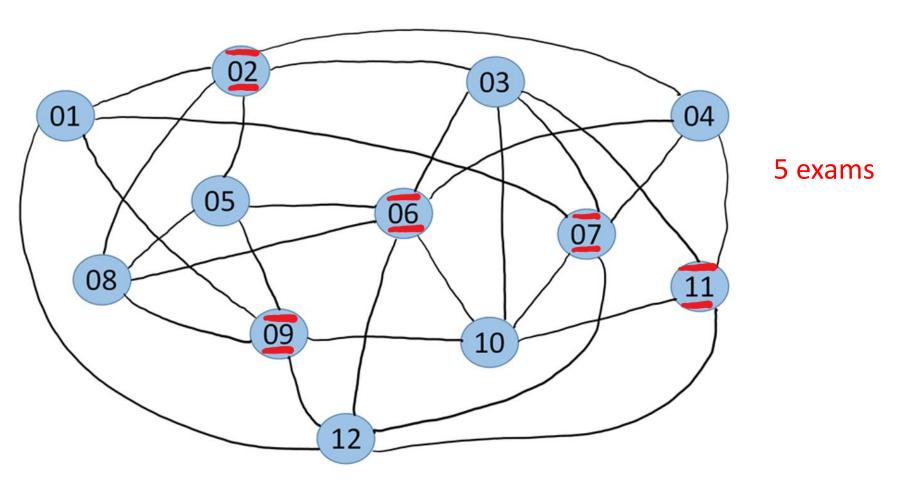
Can you do better?





■Can you do better?





- ■Tired of trying?
- ■Is this the best possible?



## 5 is the best possible!

### Proof by magic?



- ■Unfortunately not ⊗
- ■No magician in the world has pulled out such a rabbit to this day! (By this we mean a trick that would work on *all* graphs.)

- Of course there is always a proof:
  - ■Try all possible subsets of 6 nodes.
  - ■There are 924 of them.
  - Observe that none of them work.
- ■But this is no magic. It impresses nobody. We want a "short" proof. (We will formalize what this means.) Like the one in our max-flow example.
- Let's appreciate this further...



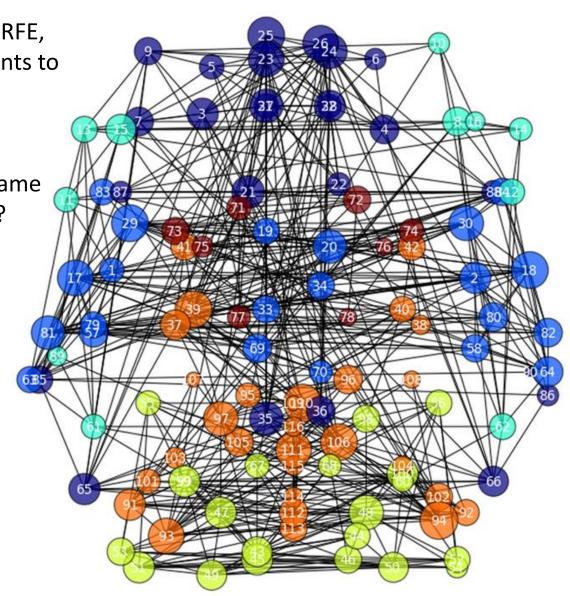
# Let's try another graph

■Encouraged by the success of ORFE, now the Dean of Engineering wants to the same for 115 SEAS courses.

■How many final exams on the same day are possible? Can you do 17?

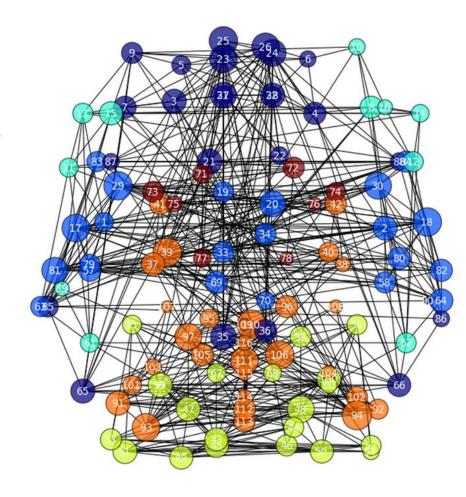
You have 6 minutes!;)

- ■Want to try out all possibilities for 17 exams?



### But there is some good news

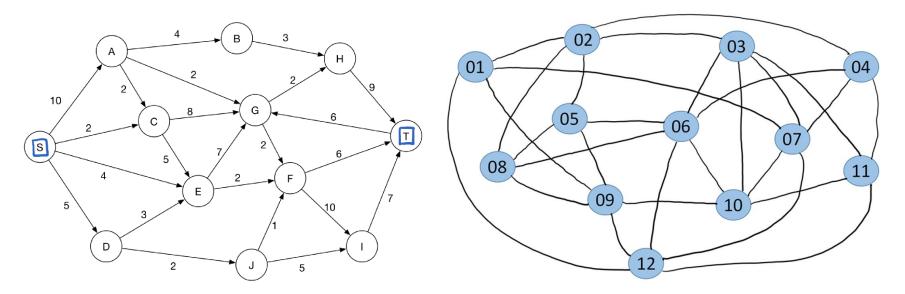
- ■Even though finding the best solution always may be too much to hope for, techniques from optimization (and in particular from the area of *convex optimization*) often allow us to find high-quality solutions with performance guarantees.
- ■For example, an optimization algorithm may quickly find 16 concurrent exams for you.
- ■You really want to know if 17 is impossible. Instead, another optimization algorithm (or sometimes the same one) tells you that 19 is impossible.



- ■This is very useful information! You know you got 16, and no one can do better than 19.
  - ■We sill see a lot of convex optimization in this class!



# Which of the two problems was harder for you?



- ■Not always obvious. A lot of research in optimization and computer science goes into distinguishing the "tractable" problems from the "intractable" ones.
- ■The two brain teasers actually just gave you a taste of the P vs. NP problem. (If you have not heard about this, that's OK. You will soon.)
- ■The first problem we can solve efficiently (in "polynomial time").
- ■The second problem: no one knows. If you do, you literally get \$1M!
  - More importantly, your algorithm immediately translates to an efficient algorithm for thousands of other problems no one knows how to solve.

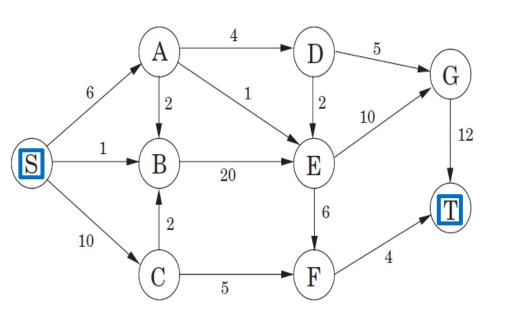




# Modelling problems as a mathematical program



### Let's revisit our first game

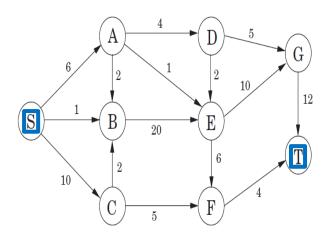


- •What were your decision variables?
- •What were your constraints?
- •What was your objective function?

### Rules of the game:

- Cannot exceed capacity on the edges.
- ■For each node, except for S ant T, flow in = flow out (i.e., no storage).
- ■Goal: ship as much oil as you can from S to T.





$$\chi_{s_A}$$
,  $\chi_{AD}$ ,  $\chi_{BE}$ , ...,  $\chi_{GT}$  Decision variables

max. 
$$\chi_{s_A} + \chi_{s_B} + \chi_{s_C}$$
 Objective function

$$\alpha$$
  $\chi_{s_A}$ ,  $\chi_{AD}$ ,  $\chi_{BE}$ , ...,  $\chi_{GT}$ 

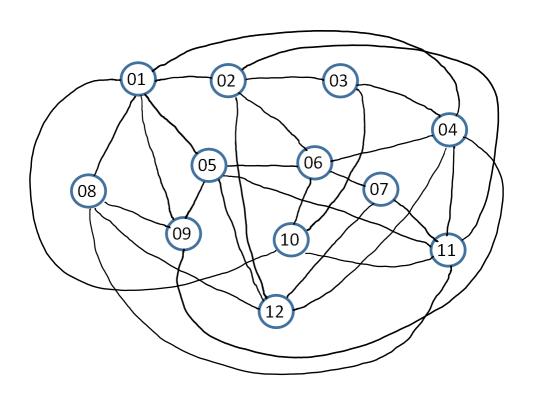
s.t.

o 
$$\mathcal{X}_{SA}$$
,  $\mathcal{X}_{AD}$ ,  $\mathcal{X}_{BE}$ , ...,  $\mathcal{X}_{GT}$  > o  $\mathcal{X}_{SA} \leqslant 6$ ,  $\mathcal{X}_{AB} \leqslant 2$ ,  $\mathcal{X}_{EG} \leqslant 10$ , ...,  $\mathcal{X}_{GT} \leqslant 12$ 

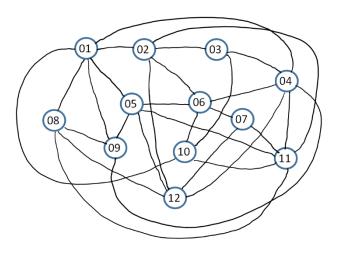


Constraints

### Let's revisit our second game



- •What were your decision variables?
- ■What were your constraints?
- •What was your objective function?



$$\chi_1, \chi_2, \ldots, \chi_{Q}$$
 Decision variables

$$\max_{x_1 + x_2 + \cdots + x_{12}}$$
 Objective function s.t.

o 
$$\chi_{i}(1-\chi_{i})=0$$
,  $i=1,...,12$ 

Constraints

o  $\chi_{i+\chi_{2}\leq 1}$ 
 $\chi_{i+\chi_{3}\leq 1}$  (one per edge)

 $\chi_{4+\chi_{4}\leq 1}$ 
 $\chi_{12+\chi_{3}\leq 1}$ 



# Why one hard and one easy? How can you tell?

$$\chi_{s_A}$$
,  $\chi_{AD}$ ,  $\chi_{BE}$ , ...,  $\chi_{G\tau}$ 

max. 
$$\chi_{sA} + \chi_{sB} + \chi_{sC}$$
s.t.

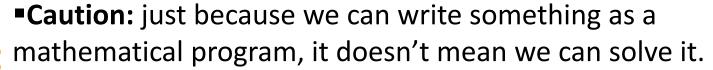
$$\circ \quad \chi_{s_A}, \ \chi_{AD}, \chi_{BE}, \ldots, \chi_{G\tau} >\!\!> \circ$$

o 
$$\chi_{SA} \leq 6$$
,  $\chi_{AB} \leq 2$ ,  $\chi_{EG} \leq 10$ ,...,  $\chi_{GT} \leq 12$ 

$$\chi_1, \chi_2, \ldots, \chi_{l_2}$$

$$\max$$
.  $\chi_1 + \chi_{1+} - \cdots + \chi_{12}$ 
s.t.

o 
$$\chi_{i}(1-\chi_{i})=0$$
,  $i=1,...,12$ 





### **Fermat's Last Theorem**

■Can you give me three positive integers x, y, z such that

$$x^2 + y^2 = z^2$$
?

And there are infinitely many more...

■How about 
$$x^3 + y^3 = z^3$$
?

■How about 
$$x^4 + y^4 = z^4$$
?

■How about 
$$x^5 + y^5 = z^5$$
?

### Fermat's Last Theorem

### Fermat's conjecture (1637):

For  $n \geq 3$ , the equation  $x^n + y^n = z^n$  has no solution over positive integers.

Proved in 1994 (357 years later!) by Andrew Wiles.

(Was on the faculty in our math department until a few years ago.)





### Arithmeticorum Liber II.

internallum numerorum 2. minor autem c' iréc. ò dea unilon isus c' irèc u' B. Din-1 N. atque ideo maior 1 N. + 2. Oportet itaque 4 N. + 4. triplos effe ad 2. & adhuc superaddere 10. Ter igitur 2. adscitis vnitatibus 10. æquatur 4 N. + 4. & fit 1 N. 3. Erit ergo minor 3. maior 5. & farisfaciunt quattioni.

σει άρα άριθμος δ' μοιάδας δ' τριπλασίετας i B. C its varpixen wir topic des Manage & his to 1. Land eight com of Mondal d'. 2) วกาส อ ฉ่อเชียง น ว . เรอเอ แล เมตรour u' 7. o d'e millon m' i. z' minst re

### IN QVAESTIONEM VIL

CONDITIONIS appositat eadem rattio est qua & appositat pracedenti quastioni, nil enim Canones iidem hie etiam locum habebunt, ve manifestum est.

### QVÆSTIO VIII.

PROPOSITY M. quadratum diuidere induos quadratos, Imperatum fit ve 16. dinidatur in duos quadratos, Ponatur primus 1 Q.Oportet igitur 16-1 Q.æquales effe quadrato. Fingo quadratum à numeris quotquot libuerit, cum defectu tot vnitatum quod continet latus ipfius 16. esto à 2 N.-4. ipie igitur quadratus eric 4Q. + 16. -16 N. hac aquabuntur vnitatibus 16 -1 Q. Communis adiiciatur vtrimque defectus, & à fimilibus auferantur fimilia, fient 5 Q. æquales 16 N. & fit 1 N. # Erit igitur alter quadratorum #. alter verò # & vtriufque fumma eft fr feu 16. & vterque quadratus eft.

TON Straybir a retpayarer d'inheir eic dús respayanous. invertaglos d'a t es διελείν είς δύο τέτρας ώτους. καλ τετάχθω ο megros duvaperes puas. deban dea ponddas is heifer d'oranemes mas Tras in The τομίων. πλάωτω τ τετράγωνον λαθ εξ. δουν dh more deifer roourur u' bour beir i ? is " made is is to to B solfe u' f. airde מפת ל דוקמים זיכנו לעשמעונים ל עו וב Acidet ce is. ratra ira peram is heidet δωνάμιως μιας. κοινή προσκείδω κ λείψε. κ) λου όμωαν όμεια. δικάκεις άρα ί ίσας ล้วเป็นอัง เรื. หู ว่าเราะ 6 ล้อยชินดิง เรื. พ่นส-โลร. โรลเอ๋ เนีย ธารี เพราราพันพ์ในร. อ๋ ฮิโ เนฮ์ไ είτος οπίμετίζων , Θ οί δύο στιστιθέντες ποιέσ

v einosomunta, inu uradus if. zal im trange neagard.

### OBSERVATIO DOMINI PETRI DE FERMAT.

C'Ebum autem in duos cubes , aut quadratoquadratum in duos quadratoquadratos & generaliter nullam in infinitum yltra quadratum potestatem in duos ciusdem nominis fas eft dividere cuius rei demonstrationem mirabilem sane detexi. Hanc marginis exignitas non caperes.

R V R S V S oporteat quadratum 16 dividere in duos quadratos, Ponatur rurfus primi latus : N. alterius verò quotcunque numerorum cum defectu tot vnitatum, quot conftat latus diuidendi. Esto itaque 2 N. - 4. erunt quadrati, hic quidem 1 Q. ille vero 4 Q. + 16. - 16 N. Cæterum volo vtrumque fimul æquari vnitatibus 16. lgitur 5 Q. + 16. -16 N. aquatur vnitatibus 16. & fit 1 N. 4 erit ΕΣΤΩ δη πάλη τὸς ις περάγωτος διελείς εἰς δύο τετραγώτους. τετάρθω πάλη או דול שטעורטע מאמלפט כי ווילכ, וי ין דול ודיוףם er omerdinnere bei fer u' bouer bei ni ru daypulden midpa. Tow o'n cf B seile u. J. להרדב: כו דודנת שוסו לכ נולי לעשמנושנ עובר, be de Suvamur & u' is heifer ce is . Bu-אם נות לוב שלים אנו שני סנידום לודות ושפע בלו על is. Swame aga i u' is hite ce is tray על וה. אפן אויותו ב מפולעולה וה מינודומים



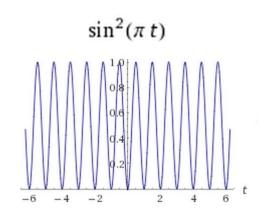
### Fermat's Last Theorem

### ■Fermat's conjecture (1637):

For  $n \ge 3$ , the equation  $x^n + y^n = z^n$  has no solution over positive integers.

■Consider the following optimization problem (mathematical program):

min. 
$$(x^{n} + y^{n} - z^{n})^{2}$$
  
s.t.  $(x^{n} + y^{n} - z^{n})^{2}$   
 $s.t.$   $(x^{n} + y^{n} - z^{n})^{2}$ 



- ■Innocent-looking optimization problem: 4 variables, 5 constraints.
- ■If you could show the optimal value is non-zero, you would prove Fermat's conjecture!



### **Course objectives**

### ■The skills I hope you acquire:

- Ability to view your own field through the lens of optimization and computation
  - ■To help you, we'll draw applications from operations research, statistics, finance, machine learning, engineering, ...
- Learn about several topics in scientific computing
- •More mathematical maturity and ability for rigorous reasoning
  - ■There will be some proofs in lecture. Easier ones on homework.
- Enhance your coding abilities (nothing too fancy, simple MATLAB)
  - ■There will be a coding component on every homework and on the take-home final.
  - ■You are free to use any other programming language instead (e.g., Python)
- Ability to recognize hard and easy optimization problems
- Ability to use optimization software
  - ■Understand the algorithms behind the software for some easier subclass of problems.



## Software you need to download

Right away:

**MATLAB** 

Available for free to Princeton students

You are free to use any other programming language instead (e.g., Python). The solutions that we provide will be in MATLAB (though, I am asking the Als to add Python solutions).

■In the next week or two (will appear on HW#2 or #3):

**CVX** 

http://cvxr.com/cvx/

(If you are comfortable with Python, you are free to use CVXPY instead: https://www.cvxpy.org/)



## **Course logistics**

- See syllabus.
- Course website:

http://aaa.princeton.edu/orf363

- For those interested:
  - Princeton Optimization Seminar
  - http://orfe.princeton.edu/events

### Image credits and references:

- [DPV08] S. Dasgupta, C. Papadimitriou, and U. Vazirani. Algorithms.
   McGraw Hill, 2008.
- [Sch05] A. Schrijver. On the history of combinatorial optimization (till 1960). In "Handbook of Discrete Optimization", Elsevier, 2005. http://homepages.cwi.nl/~lex/files/histco.pdf

