

Part 2: Planning and Analysis Tools of Transportation Demand and Investment

Traffic Assignment:

- Constant link costs (linear networks):
 - “All-or-nothing” assignments
 - Multiple path assignments:
 - K-shortest paths
 - “Essentially-equal” shortest path
- Volume-dependent link costs
 - $C_Q = C_0 \{ 1 + a (Q / Q_{\max})^b \}$
where
 C_Q = link cost at traffic flow q
 C_0 = “zero flow” link cost
 Q = traffic flow (veh/hr.)
 Q_{\max} = practical capacity
 a, b are parameters
- System-optimum assignments
- User-optimum assignments
 - Wardrop’s first principle : There exists no path that has a lower cost.

Analytical Representations of Spatial Elements:

- Points (Nodes)
 - Location in “n”- dimensional space
 - Often 2-D; ex. Latitude, Longitude in some defined coordinate system such as [WGS-84 ellipsoid](#) [Overview](#) [Detailed equations](#) [Transformations to other datum](#) [National Geodetic Survey \(2cm accuracy objective\)](#)
 - Other attributes of nodes: name, type, number
- Lines (Arcs (Links))
 - Defined by two (2) end points (“A”-node, “B”-node)
 - Attributes:
 - Distance: [UTM Coord System](#) (using a “ruler”)
 - Name, speed limit, travel time, travel time distribution, ...
- Polygons (Surface areas (often planar))
 - Defined by ordered closed sequence of directed arcs (“Closed Directed Walk”)
 - Attributes: type, name, value,...
 - [GeoTIFF](#) (geographic referencing of TIFF file types)
- Volumes (Closed ensemble of surface areas)
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Networks:

Definitions and notation:

Directed graphs and Networks: A directed graph, $G = (N, A)$ consists of a set of N Nodes and a set of A arcs whose elements are ordered pairs of distinct nodes. A *directed network* is a directed graph whose nodes and/or arcs have associated numerical values.

Undirected graphs and networks: A directed graph with arcs having unordered pairs of distinct nodes.

Tails and Heads: A directed arc (i,j) has two end points i and j . We refer to i (“A” node) as the tail node and j (“B” node) as the head node. Arc (i,j) *emanates* from i and *terminates* at j . Arc (i,j) is incident to nodes i and j . It is *outgoing* from i and *incoming* to j . Whenever an arc $(i,j) \in A$, then node j is adjacent to i .

Degrees: The *indegree* of a node is the number on incoming arcs to that node and the *outdegree* is the number of its outgoing arcs. The *degree* of a node is the sum of the in- and out- degrees.

Adjacency list: The *arc adjacency list*, $A(i)$ of a node i is the set of arcs emanating from that node. $A(i) = \{ (i,j) \in A : j \in N \}$. The *node adjacency list*, $N(i)$ is the set of nodes adjacent to that node; $N(i) = \{ j \in N : (i,j) \in A \}$

Multiarcs and Loops: *Multiarcs* are two or more arcs having the same tail and head nodes. A *loop* is an arc whose tail node is the same as its head node.

Subgraph: A graph $G' = (N', A')$ is a *subgraph* of $G = (N, A)$ if $N' \subseteq N$ and $A' \subseteq A$. A graph $G' = (N', A')$ is a *spanning subgraph* of $G = (N, A)$ if $N' = N$ and $A' \subseteq A$.

Walk: A *walk* in a directed graph is a subgraph of G containing a connected sequence of nodes (without any mention of arcs) or a connected sequence of arcs (without any mention of nodes).

Directed Walk: A *directed walk* is an “oriented” version of a walk in that any two consecutive nodes on the walk must be a member of the set of arcs $(i_k, i_{k+1}) \in A$.

Path: A path is a walk without any repetition of nodes. We can partition the arcs of a path into two groups : forward arcs and backward arcs. An arc (i,j) in the path is a *forward arc* if the path visits node i prior to visiting node j , and is a *backward arc* otherwise.

Directed Path: A *directed path* is a directed walk without any repetition of nodes. We can store a path easily by defining a *predecessor index*, $pred(j)$ for every node j in the path.

Cycles: A *cycle* is a path $i_1 - i_2 - i_3 - \dots - i_r - i_1$. Cycles can be directed.

Acyclic Graph: A graph is *acyclic* if it contains no cycles.

Connectivity: We will say that two nodes, i and j , are *connected* if the graph contains at least one path from node i to node j . A graph is *connected* if every pair of nodes is connected; otherwise it is *disconnected*.

Strong connectivity: A connected graph is strongly connected if it contains at least one *directed* path from every node to every other node.

Cut: A cut is a partition of the node set N into two parts, S and $\underline{S} = N - S$. Each cut defines a set of arcs consisting of those arcs that have one endpoint in S and the other in \underline{S} .

$s - t$ Cut: This cut has node $s \in S$ and $t \in \underline{S}$.

Tree: A *tree* is a connected graph that contains no cycles.

A tree of n nodes contains exactly $n - 1$ arcs.

A tree has at least two *leaf* nodes (nodes with degree 1)

Every two nodes of a tree are connected by a unique path.

Forest: A graph that contains no cycles is a *forest*. Alternatively, a forest is a collection of trees.

Subtree: A connected subgraph of a tree is a *subtree*.

Rooted trees: A *rooted tree* is a tree with a specially designated node called a *root*. Rooted trees are viewed as hanging from the root.

Directed-out-tree: every path from node s is a directed path